

Finding High-Quality Local Minima in Derivative-Free Optimization

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Motivation

- ▶ We want to identify distinct, “high-quality”, local minimizers of

$$\text{minimize } f(x)$$

$$l \leq x \leq u$$

$$x \in \mathbb{R}^n$$

- ▶ High-quality can be measured by more than the objective.



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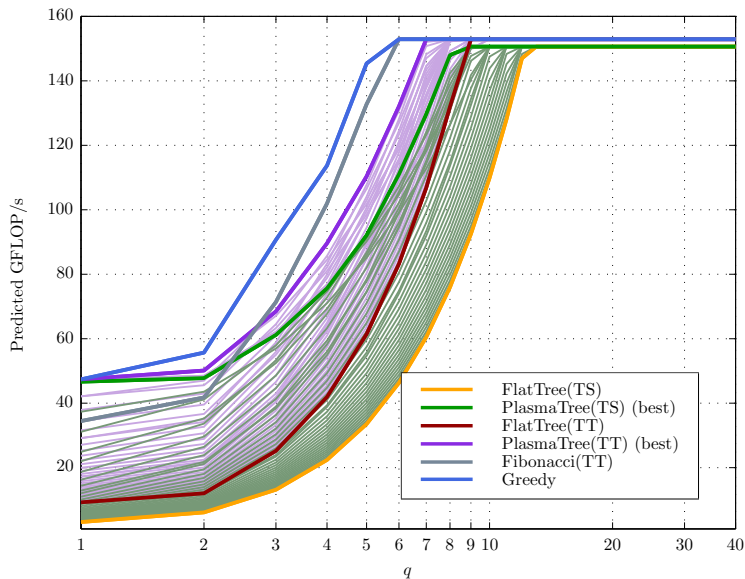
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- ▶ Derivatives of f may or may not be available.
- ▶ The simulation f is likely using parallel resources, but it does not utilize the entire machine.



Why concurrency? Tiled QR example



[Bouwmeester, et al., Tiled QR Factorization Algorithms, 2011]



Global optimization is difficult

Theorem (Törn and Žilinskas, *Global Optimization*, 1989)

An algorithm converges to the global minimum of any continuous f on a domain \mathcal{D} if and only if the algorithm generates iterates that are dense in \mathcal{D} .



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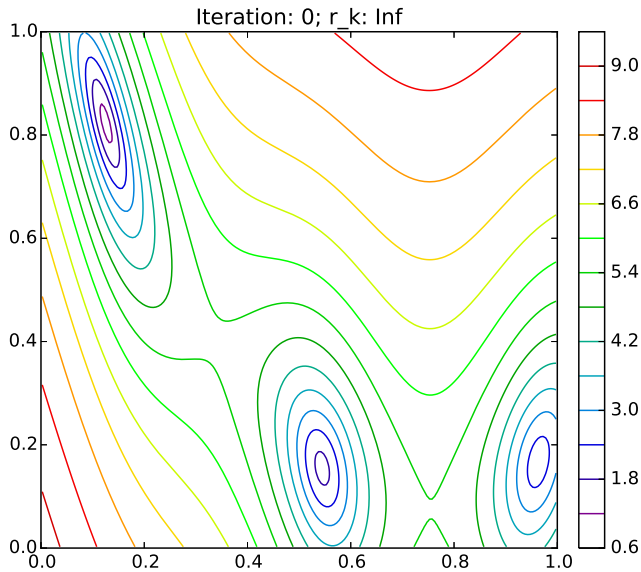
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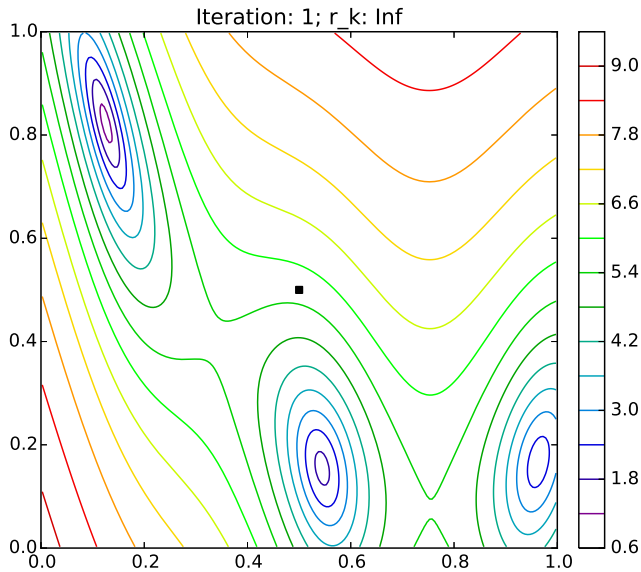
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An algorithm must trade-off between “refinement” and “exploration”.

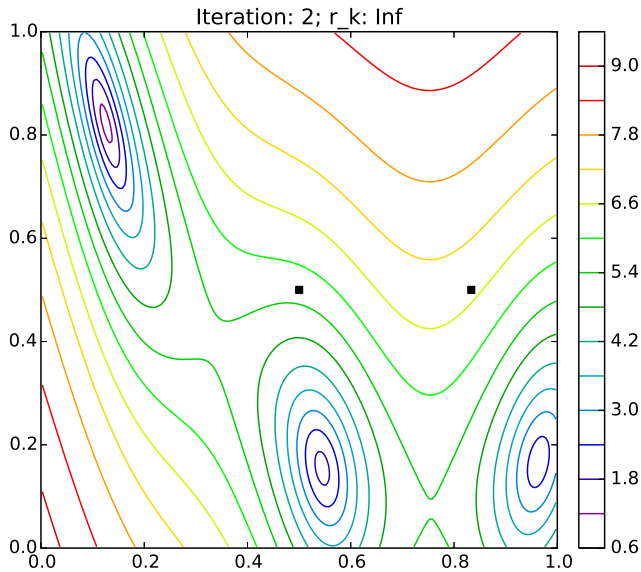
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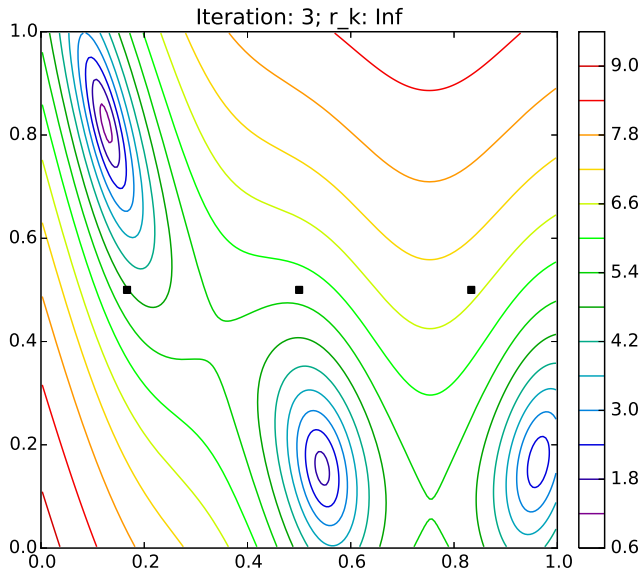
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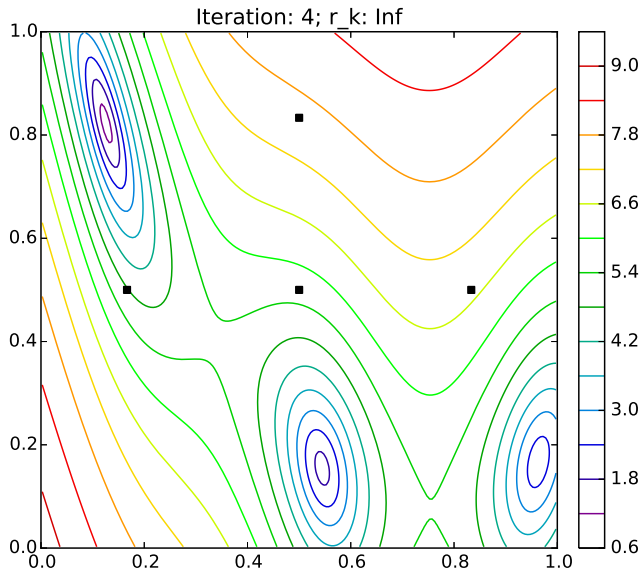
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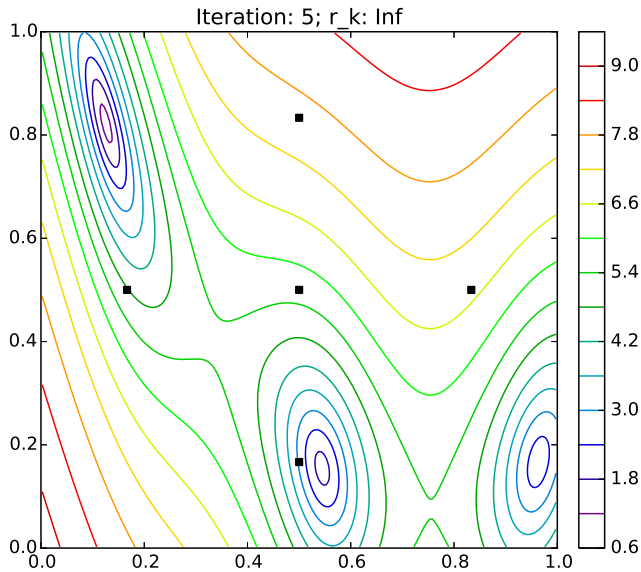
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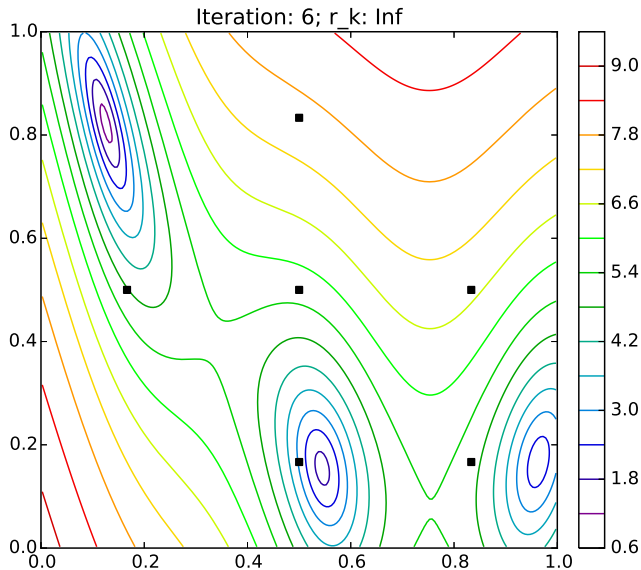
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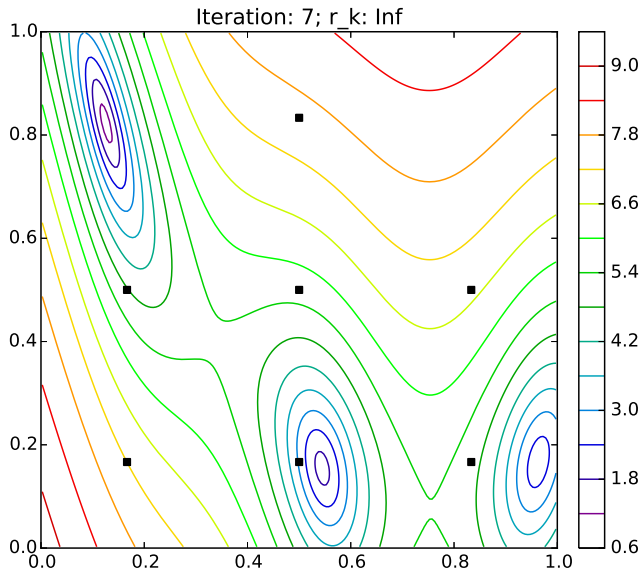
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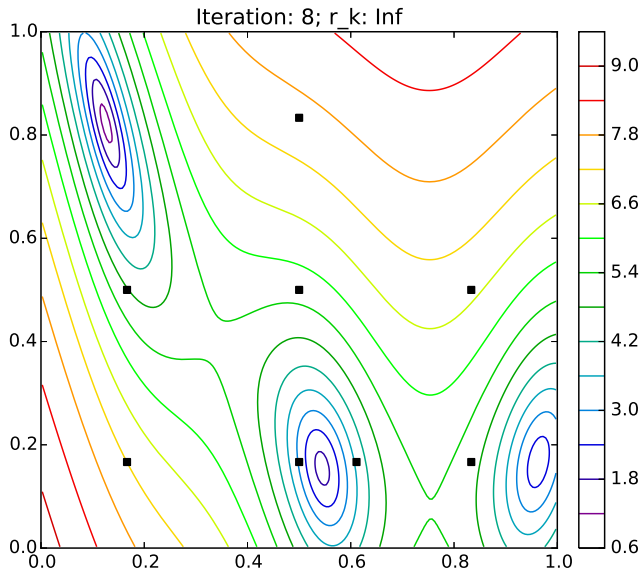
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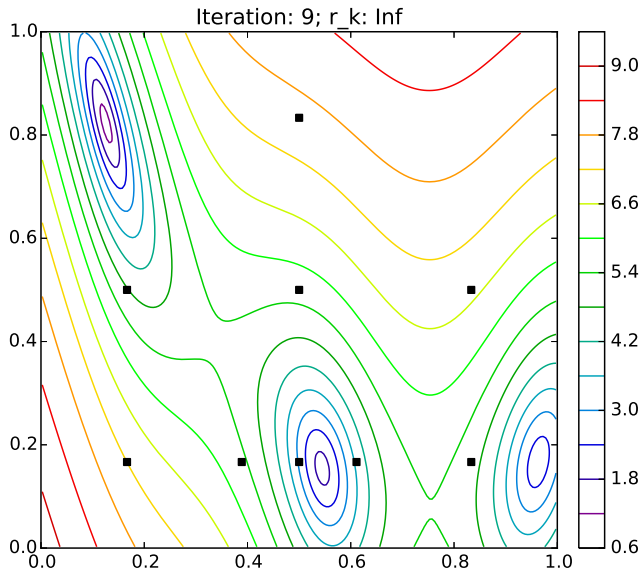
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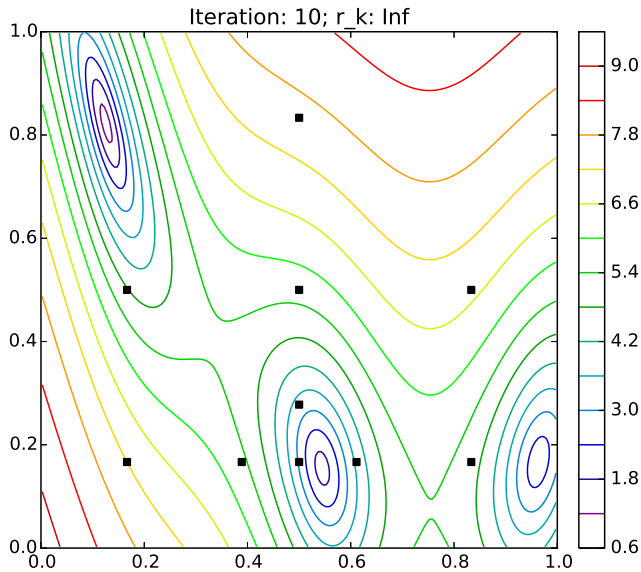
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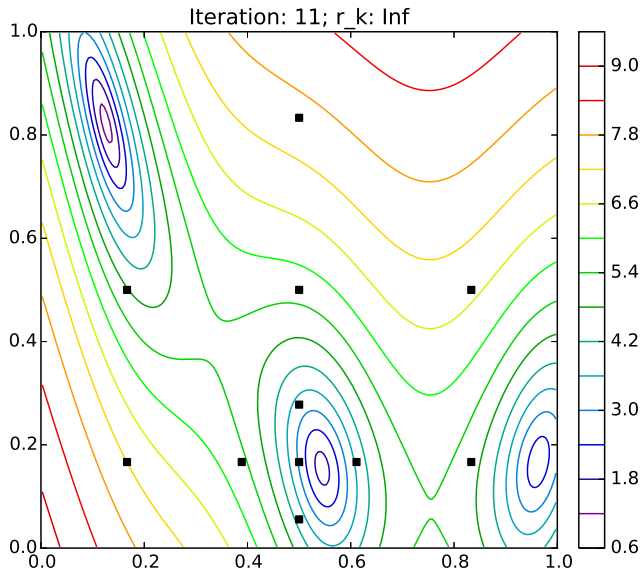
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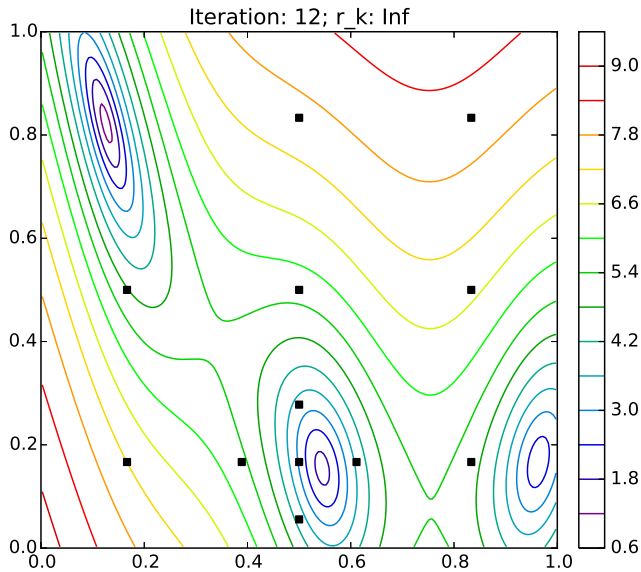
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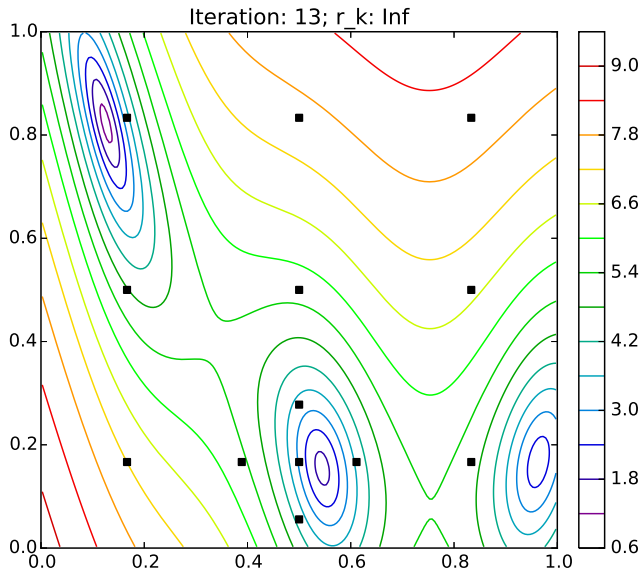
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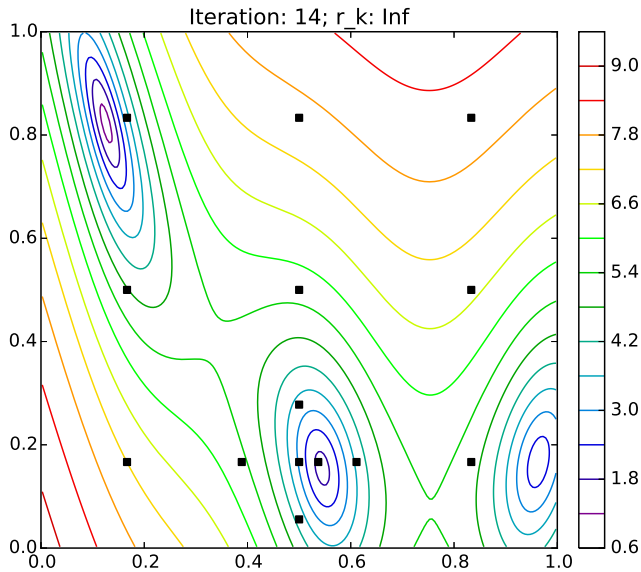
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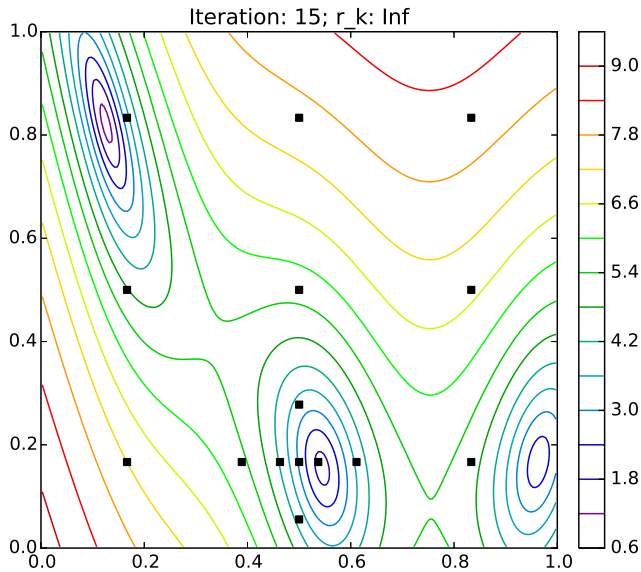
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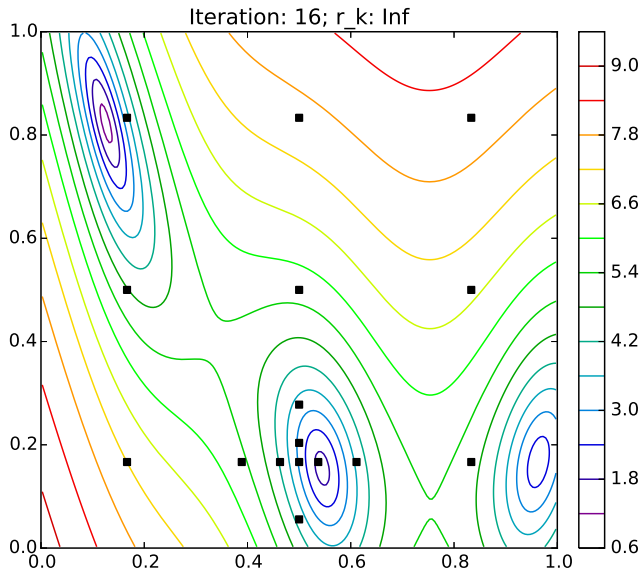
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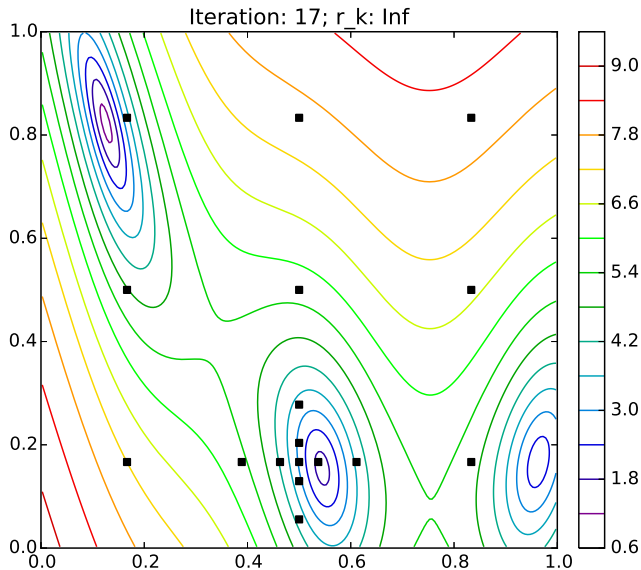
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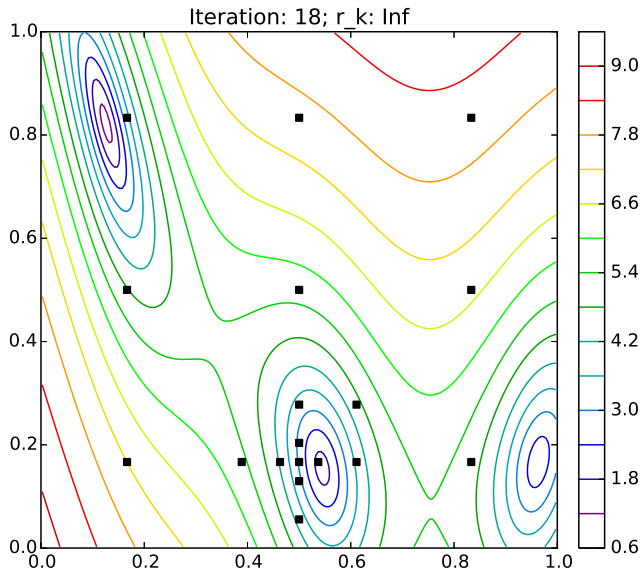
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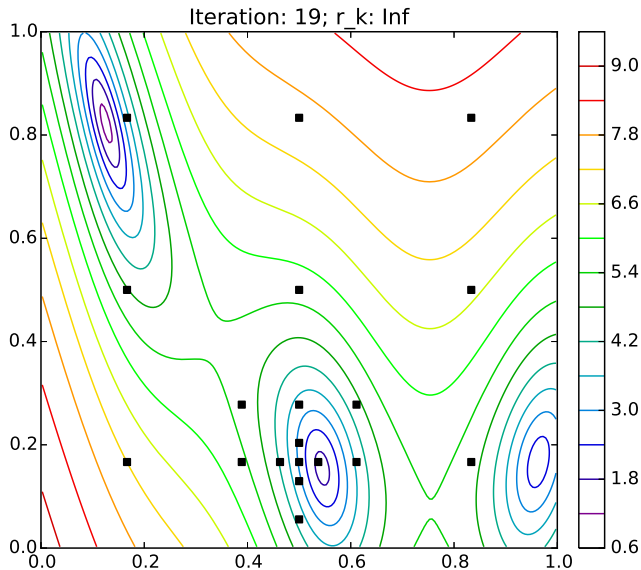
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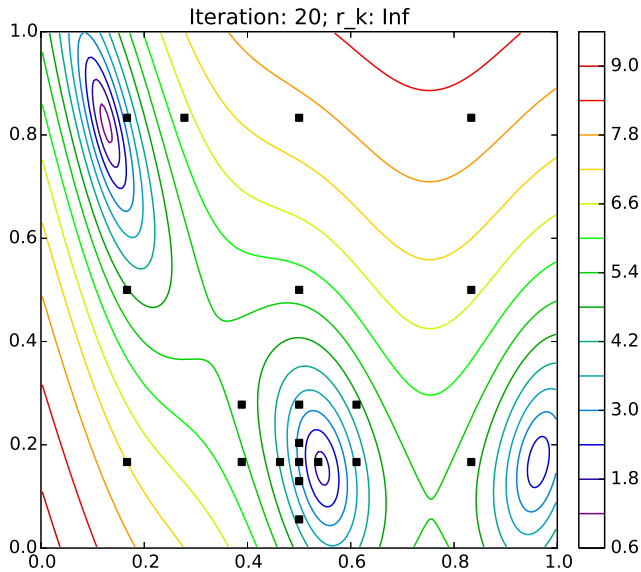
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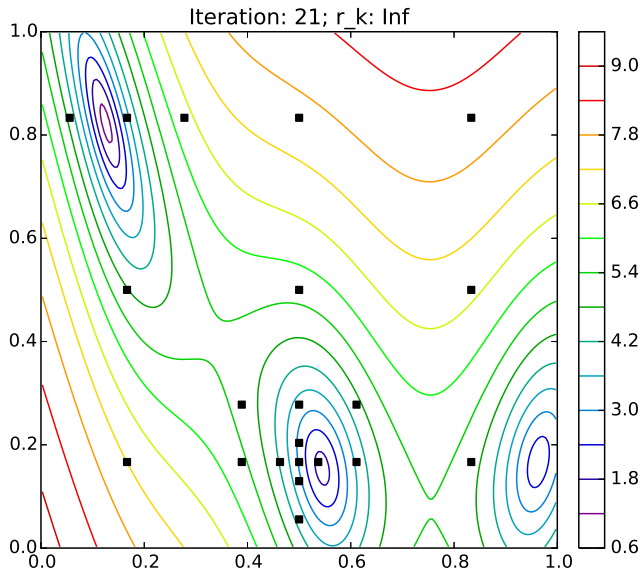
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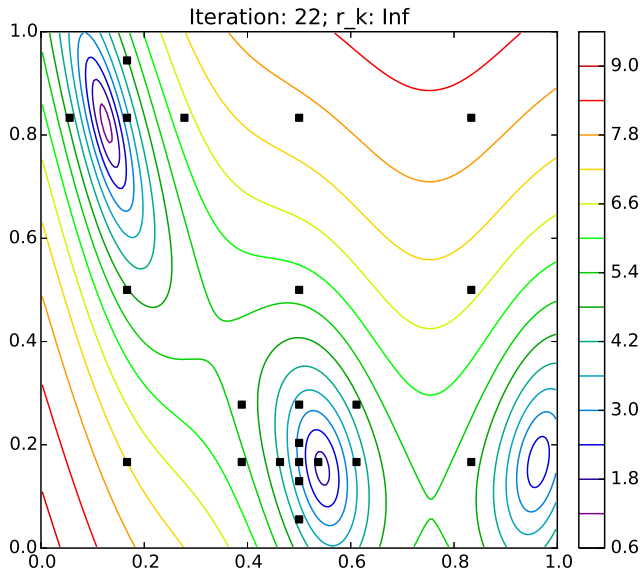
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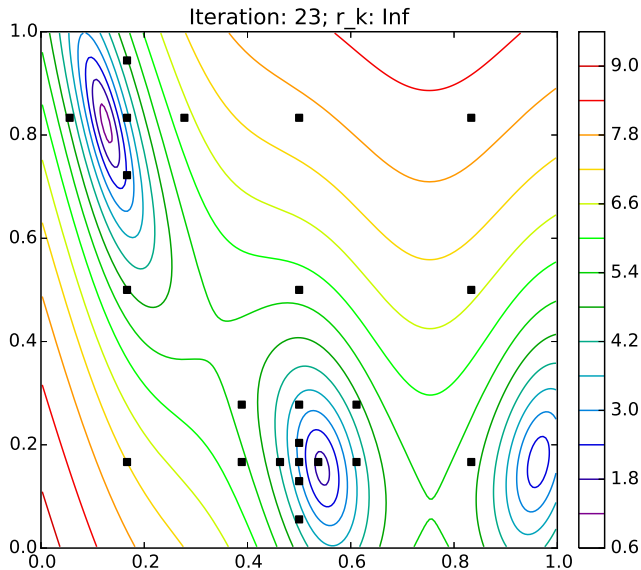
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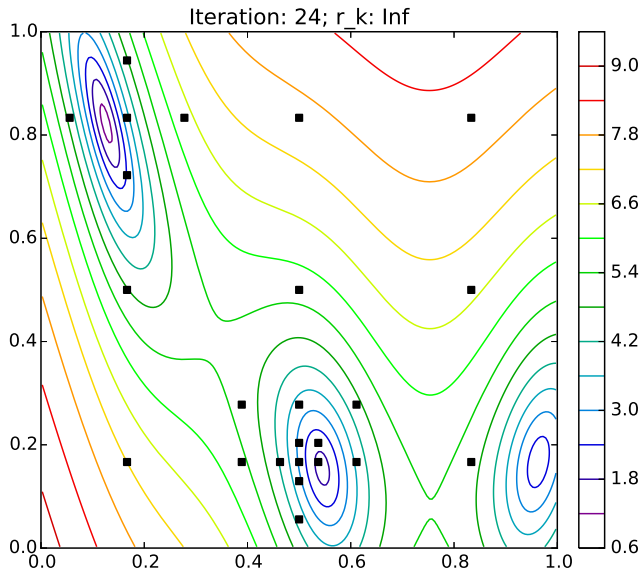
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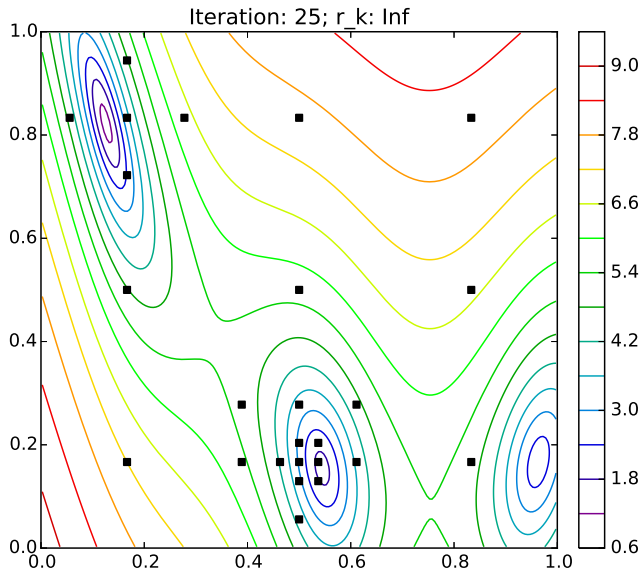
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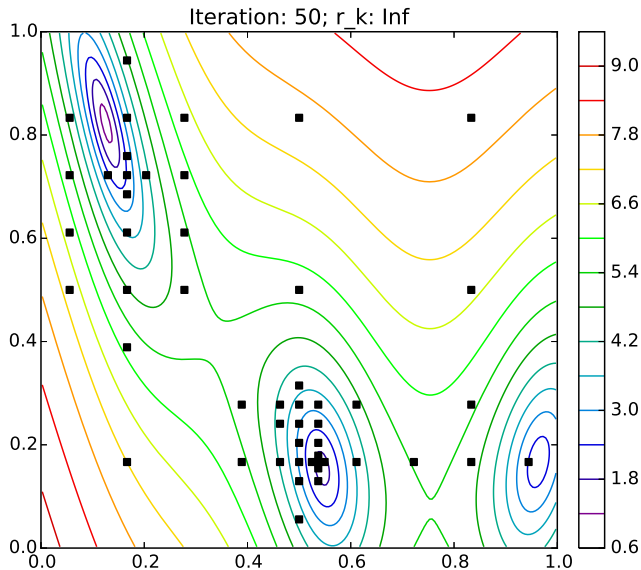
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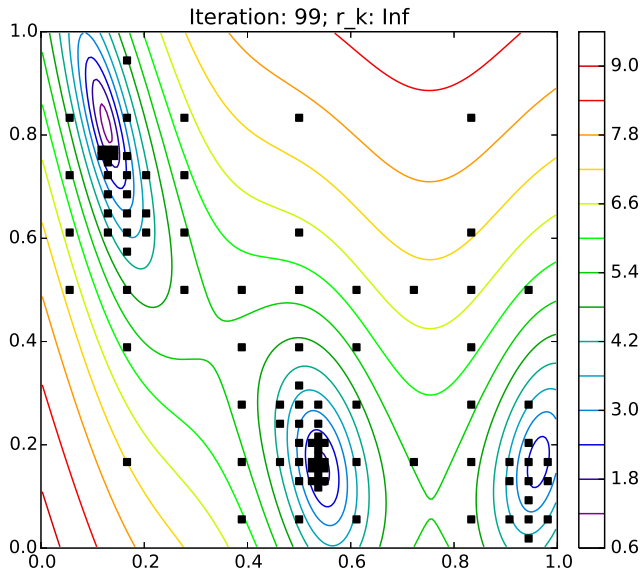
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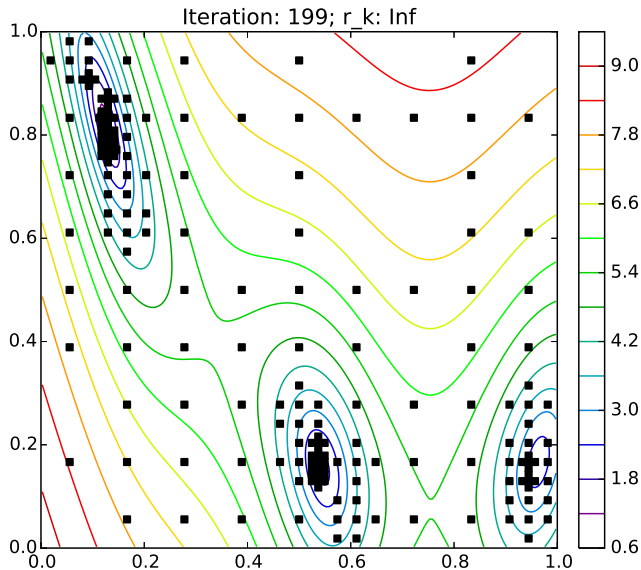
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Multistart Methods

Given some local optimization routine L :

Algorithm 1: General Multistart

for $k = 1, 2, \dots$ **do**

 Evaluate f at N points drawn from \mathcal{D}

 Start L at some set (possibly empty) of previously evaluated points



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- ▶ Possibly multiple levels of parallelism (objective, local method, global method); L may involve many sequential evaluations of $f \dots$



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- ▶ If resources are limited, how should points from each run receive priority?
- ▶ Ideally, only one run is started for each minima.
- ▶ Exploring by sampling. Refining with L .



Multi-Level Single Linkage

Given some local optimization routine L :

Algorithm 2: MLSL

for $k = 1, 2, \dots$ **do**

 Sample f at N random points drawn uniformly from \mathcal{D}

 Start L at all sample points x :

- ▶ that has yet to start a run
 - ▶ $\nexists x_i : \|x - x_i\| \leq r_k$ and $f(x_i) < f(x)$
- └
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[Rinnooy Kan and Timmer, *Mathematical Programming*, 39(1):57–78, 1987]



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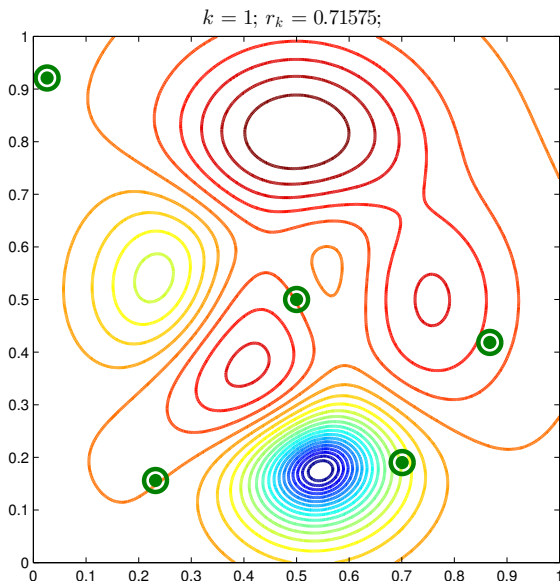
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for  $k = 1, 2, \dots$  do
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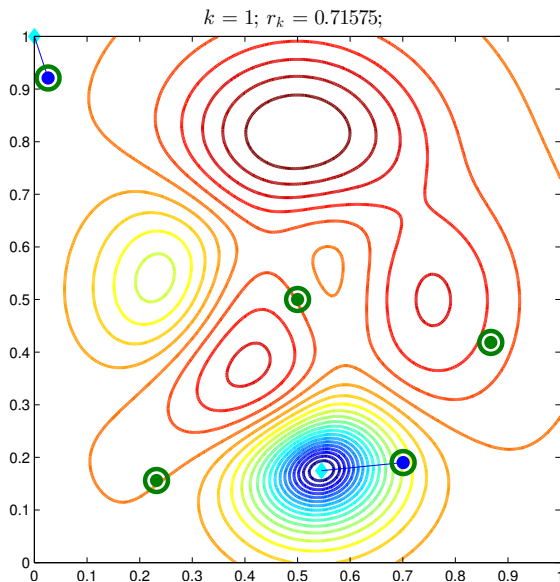
- ▶ Doesn't naturally translate when evaluations of f are limited
- ▶ Ignores some points when deciding where to start L



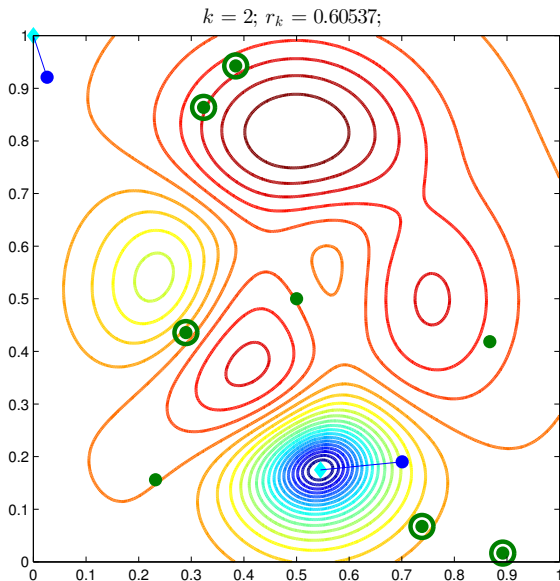
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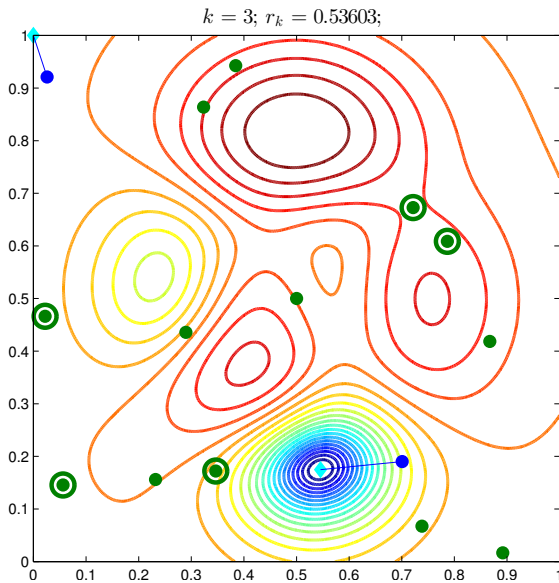
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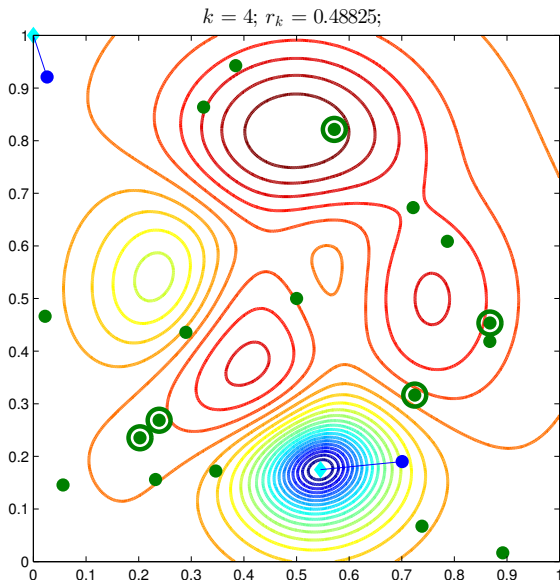
Multi-Level Single Linkage



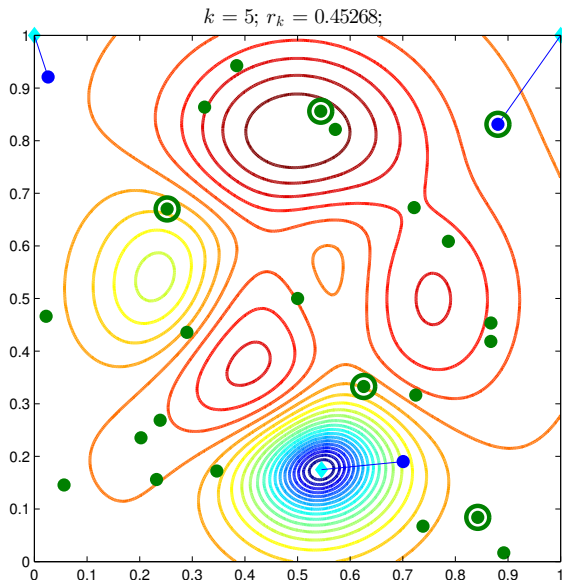
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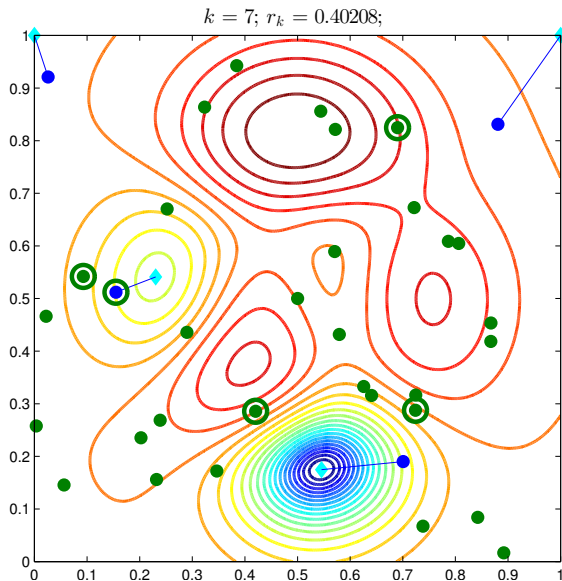
Multi-Level Single Linkage



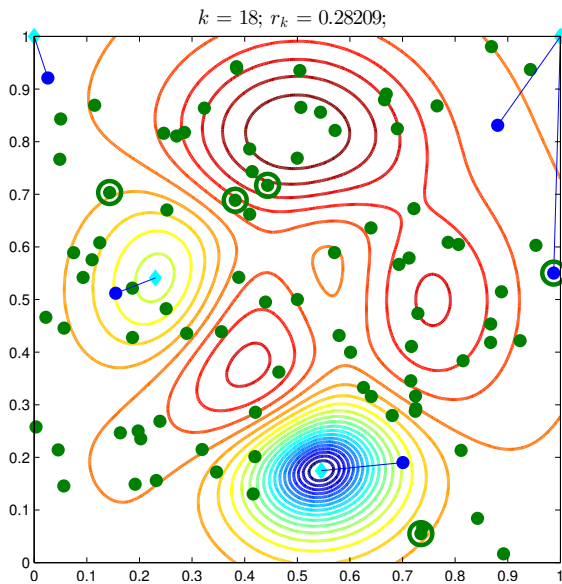
Multi-Level Single Linkage



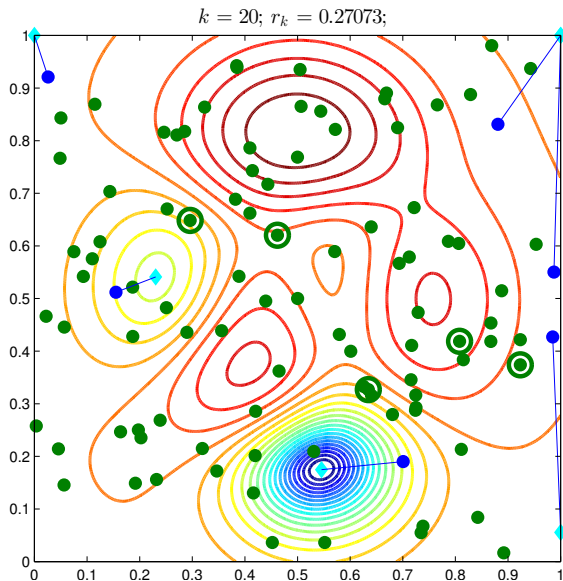
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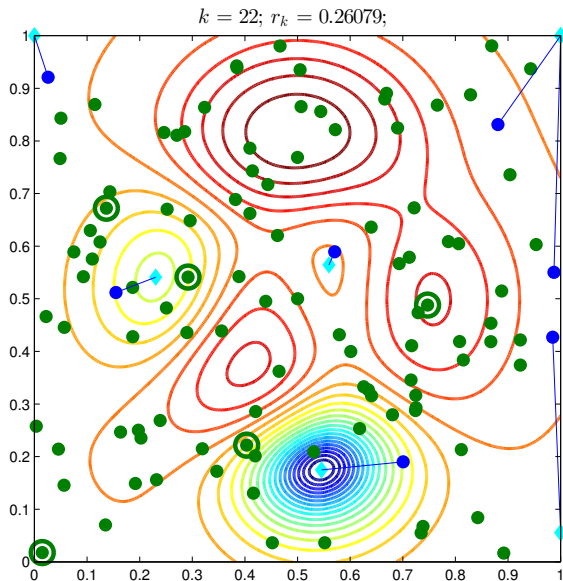
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- ▶ $f \in C^1$, with local minima in the interior of \mathcal{D} , and the distance between these minima is bounded away from zero.
- ▶ L is strictly descent and converges to a minimum (not a stationary point).

▶

$$r_k = \frac{1}{\sqrt[n]{\pi}} \sqrt[n]{\Gamma\left(1 + \frac{n}{2}\right) \text{vol}(\mathcal{D}) \frac{\sigma \log kN}{kN}} \quad (1)$$



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If r_k is defined by (1) with $\sigma > 4$, even if the sampling continues forever, the total number of local searches started is finite almost surely.



POSMM

MLSL: (S2)–(S4)

$$\hat{x} \in S_k$$

- (S2) $\nexists x \in S_k$ *with*
[$\|\hat{x} - x\| \leq r_k$ *and* $f(x) < f(\hat{x})$]
- (S3) \hat{x} *has not started a local optimization run*
- (S4) \hat{x} *is at least μ from $\partial\mathcal{D}$ and ν from known local minima*



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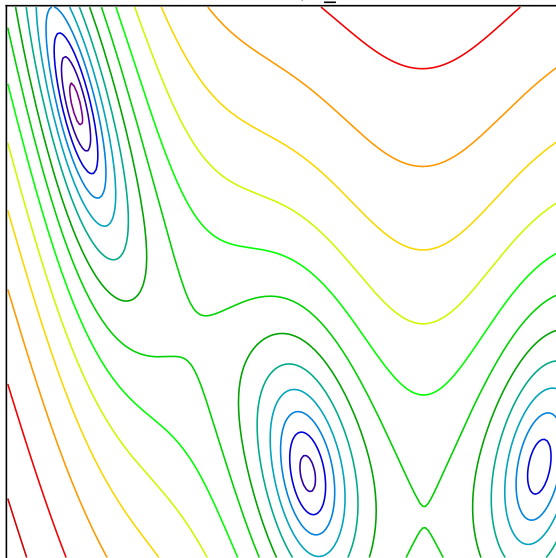
POSMM: (S1)–(S4), (L1)–(L6)

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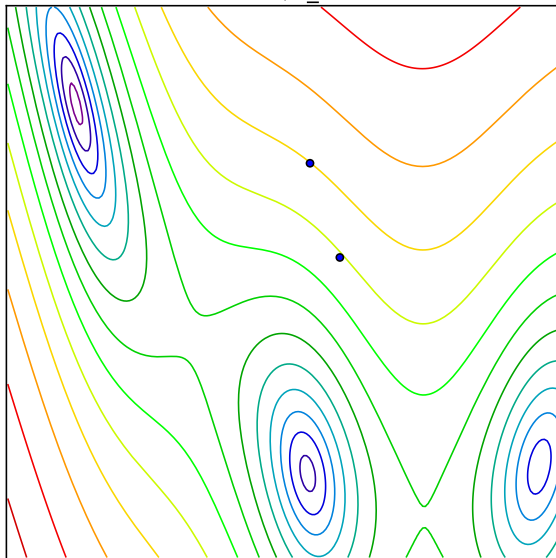
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- (L4) \hat{x} is at least μ from $\partial\mathcal{D}$ and ν
from known local minima
- (L5) \hat{x} is not in an active local
optimization run and has not
been ruled stationary
- (L6) $\exists r_k$ -descent path in H_k from
some $x \in S_k$ satisfying (S2-S4)
to \hat{x}



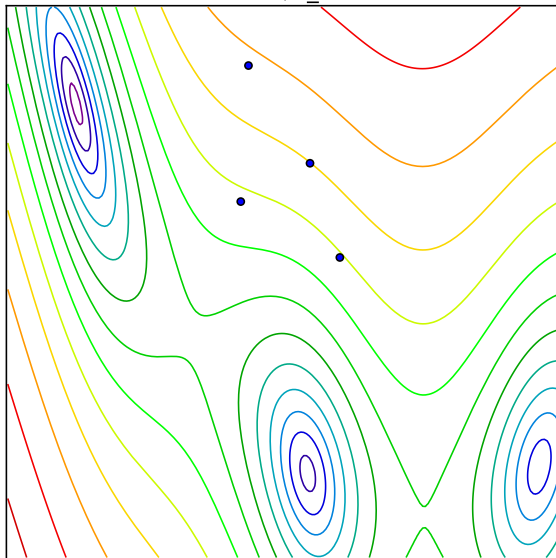
Iteration: 0; r_k : Inf



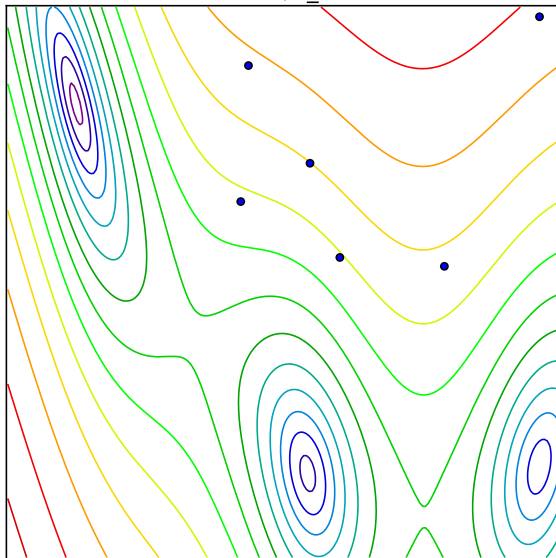
Iteration: 1; r_k : 0.743



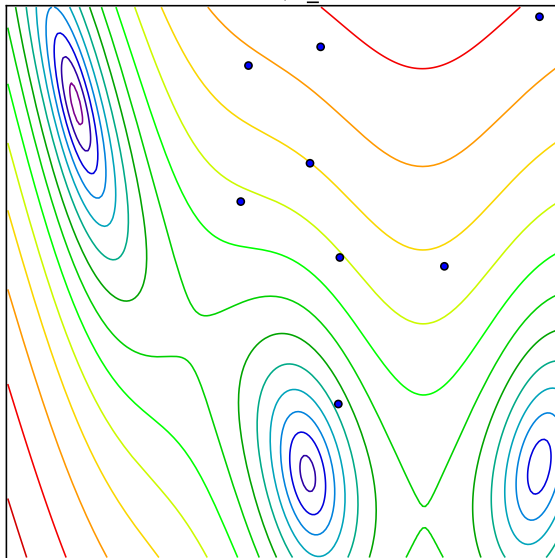
Iteration: 2; r_k : 0.743



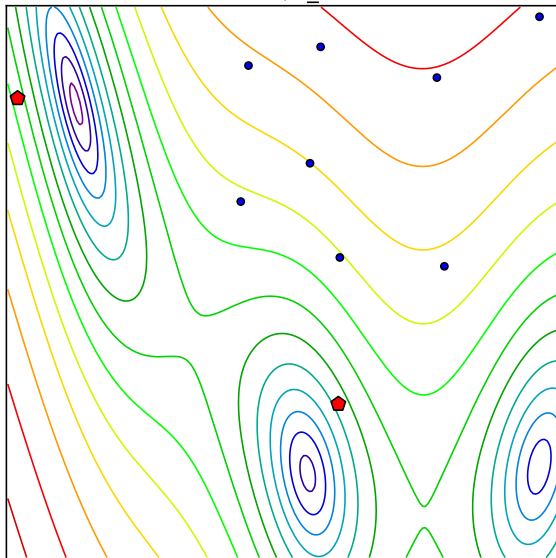
Iteration: 3; r_k : 0.689



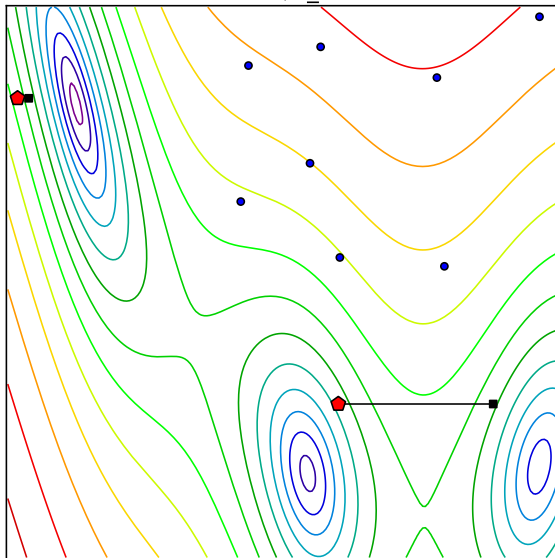
Iteration: 4; r_k : 0.643



Iteration: 5; r_k : 0.605

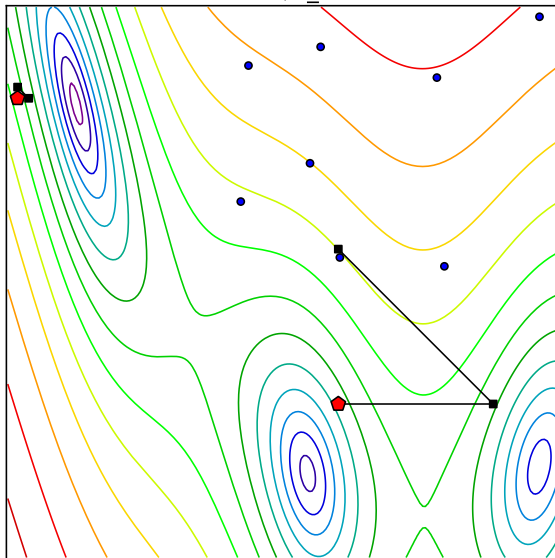


Iteration: 6; r_k : 0.605



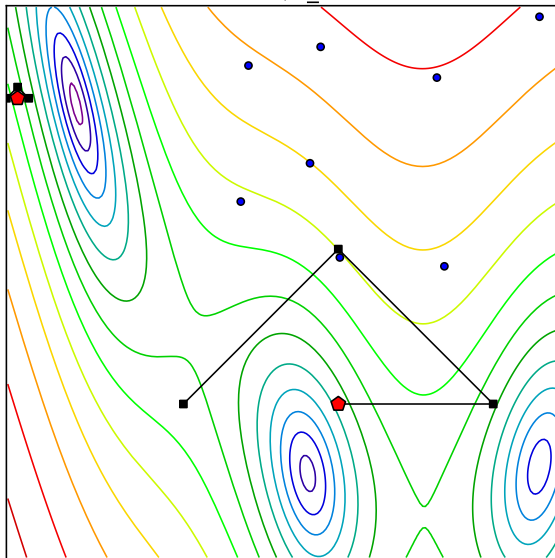
POSMM

Iteration: 7; r_k : 0.605

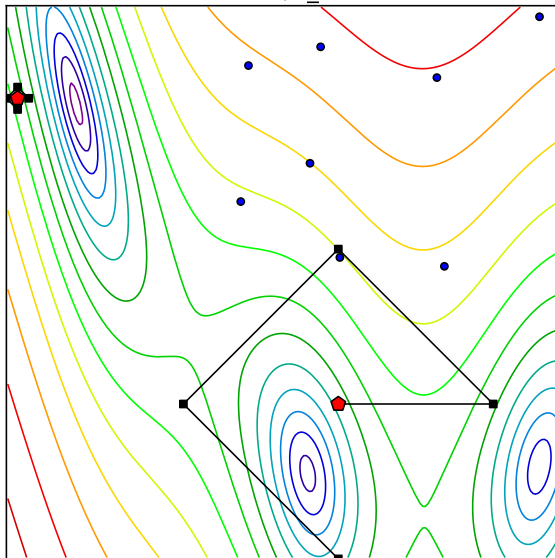


POSMM

Iteration: 8; r_k : 0.605

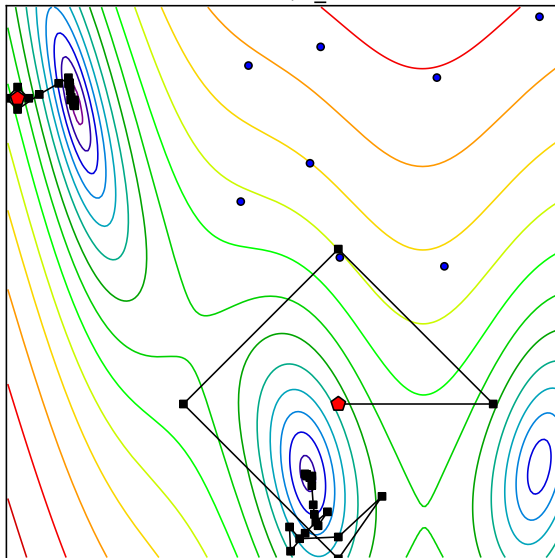


Iteration: 9; r_k : 0.605

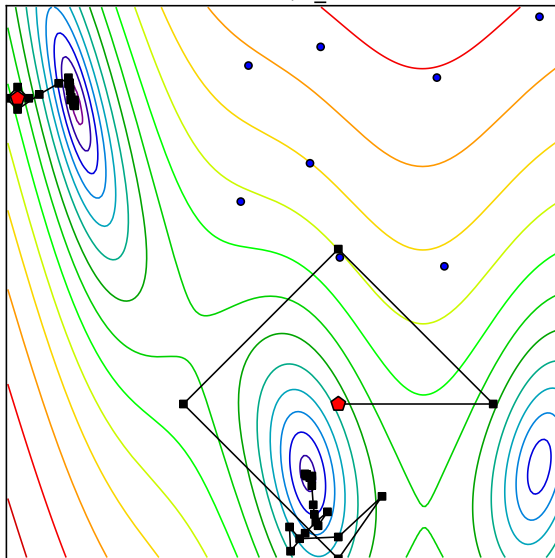




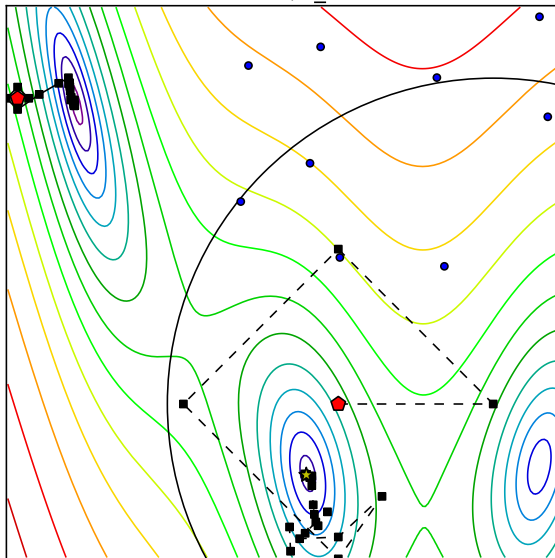
Iteration: 35; r_k : 0.605



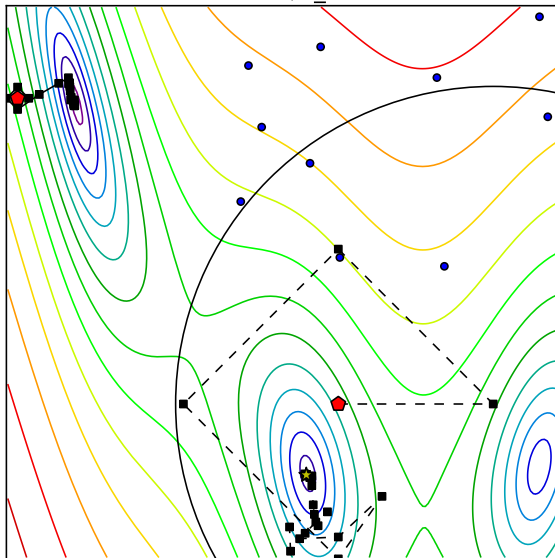
Iteration: 36; r_k : 0.605



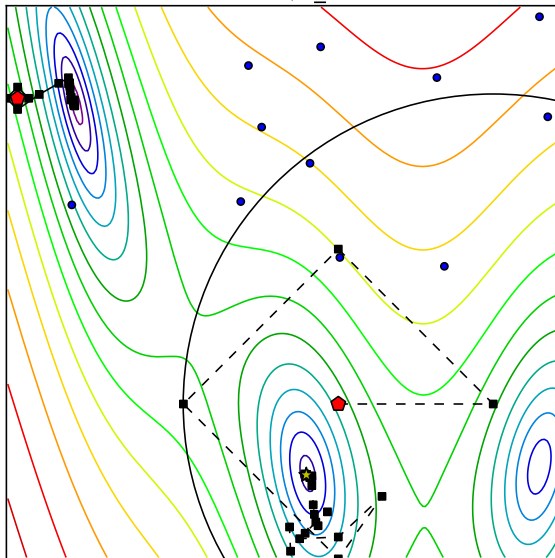
Iteration: 37; r_k : 0.589



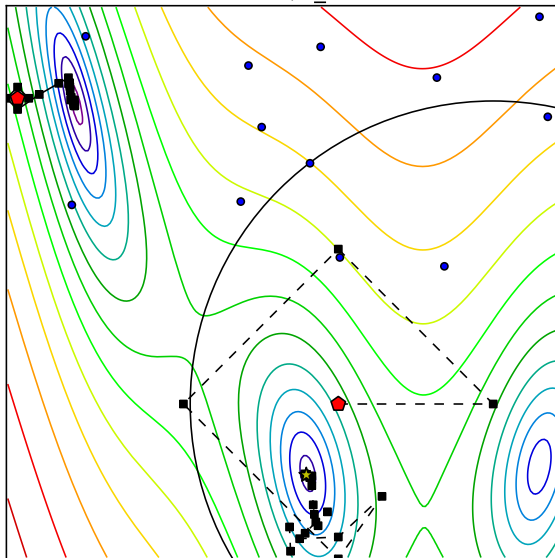
Iteration: 38; r_k : 0.574



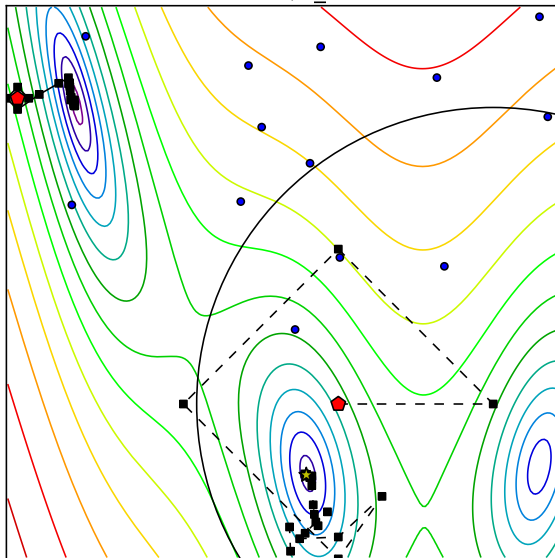
Iteration: 39; r_k : 0.560



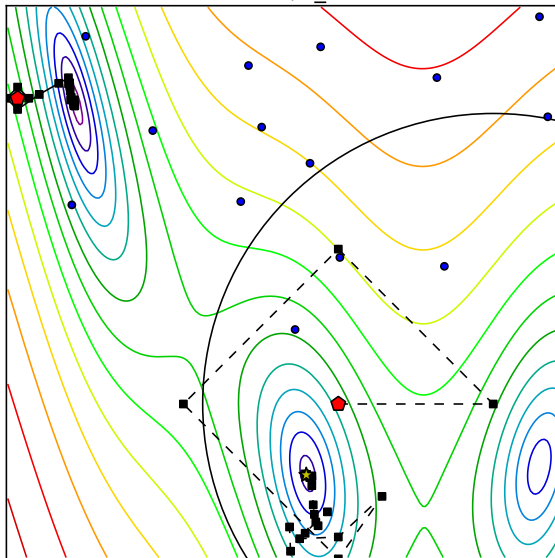
Iteration: 40; r_k : 0.548



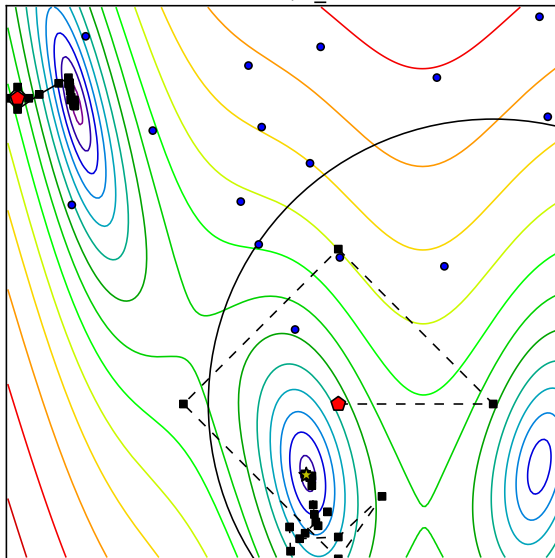
Iteration: 41; r_k : 0.536



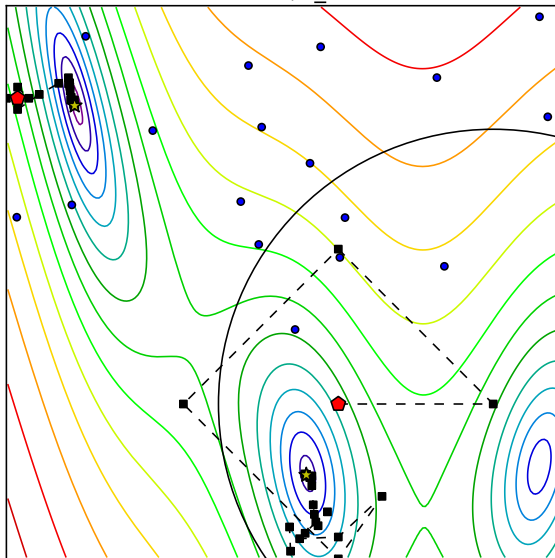
Iteration: 42; r_k : 0.525



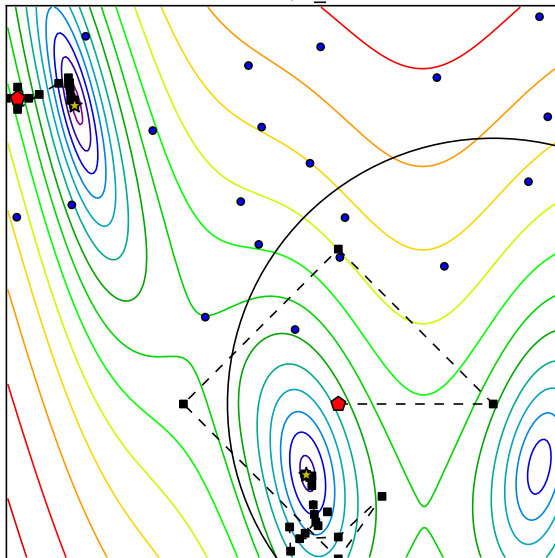
Iteration: 43; r_k : 0.515



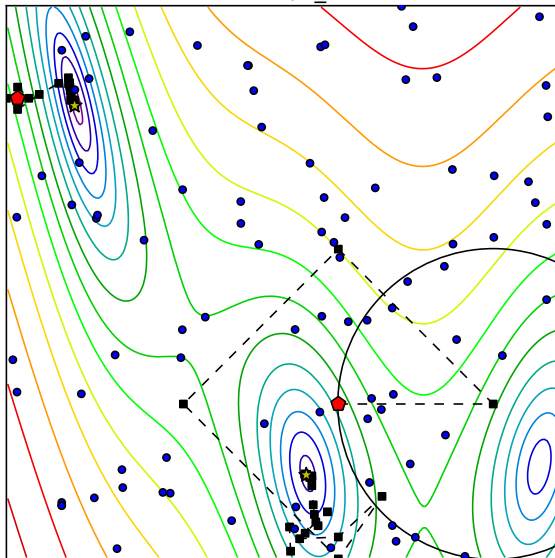
Iteration: 44; r_k : 0.497



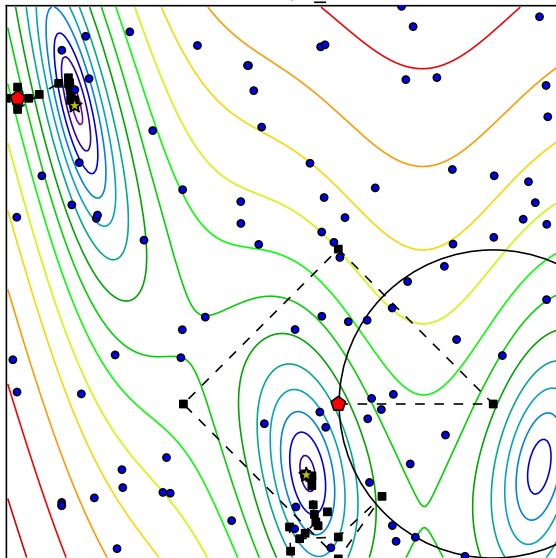
Iteration: 45; r_k : 0.480



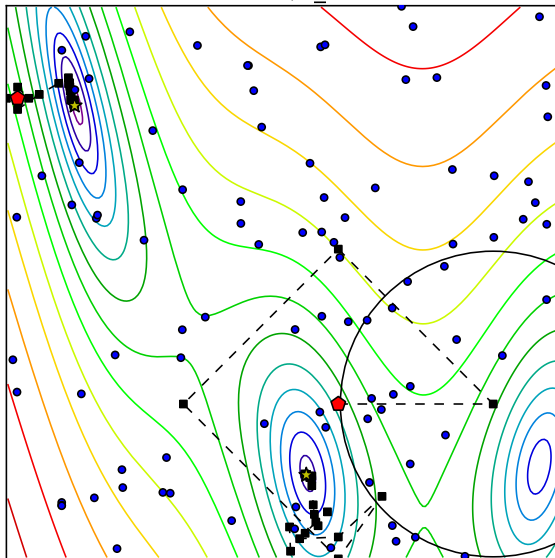
Iteration: 80; r_k : 0.281



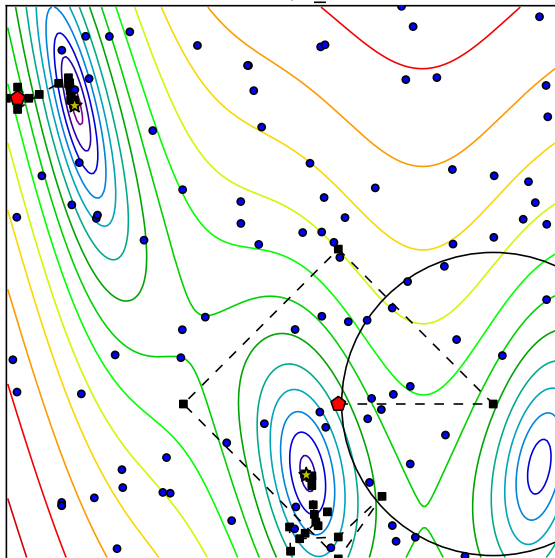
Iteration: 81; r_k: 0.279



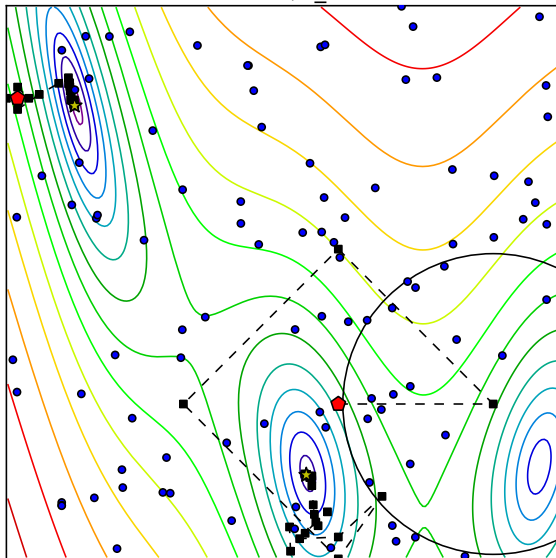
Iteration: 82; r_k : 0.276



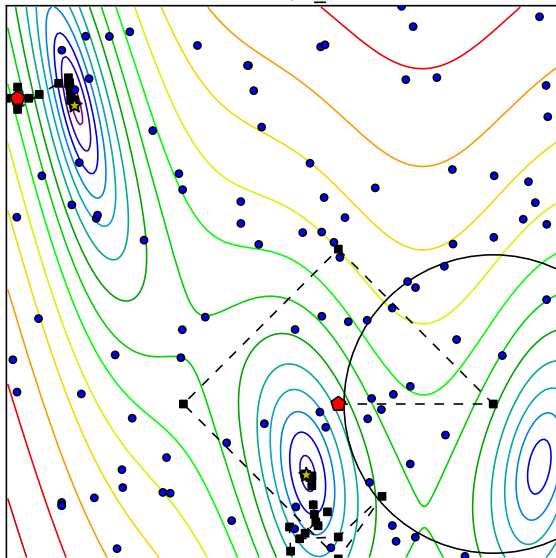
Iteration: 83; r_k : 0.274



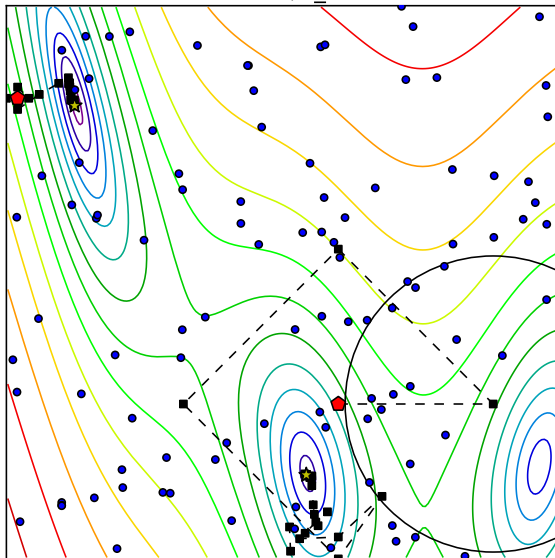
Iteration: 84; r_k : 0.272



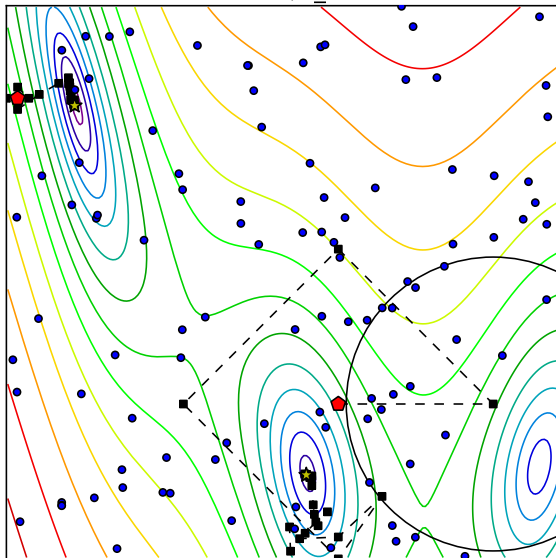
Iteration: 85; r_k : 0.270



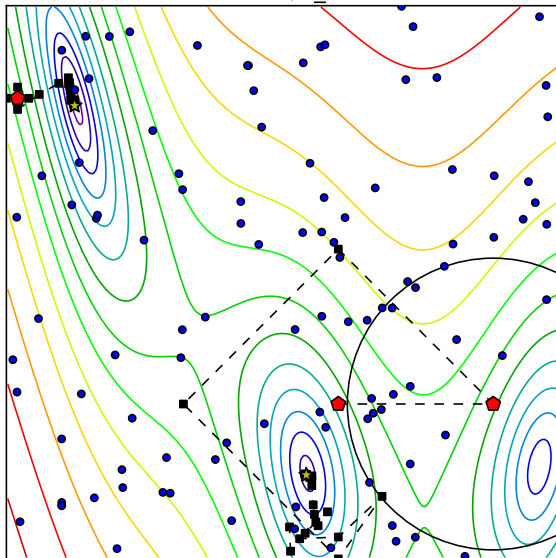
Iteration: 86; r_k : 0.268



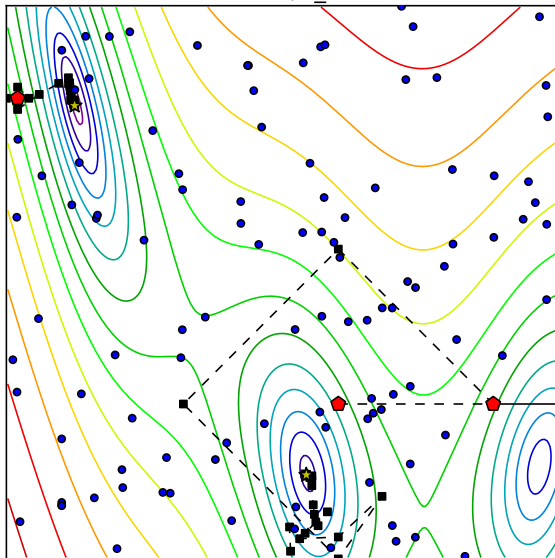
Iteration: 87; r_k : 0.266



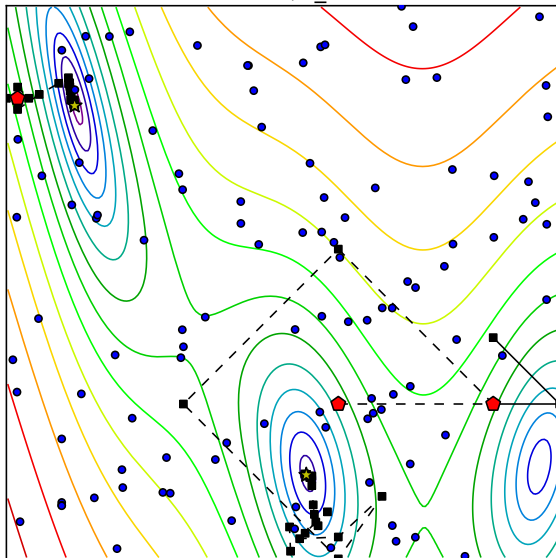
Iteration: 88; r_k : 0.264



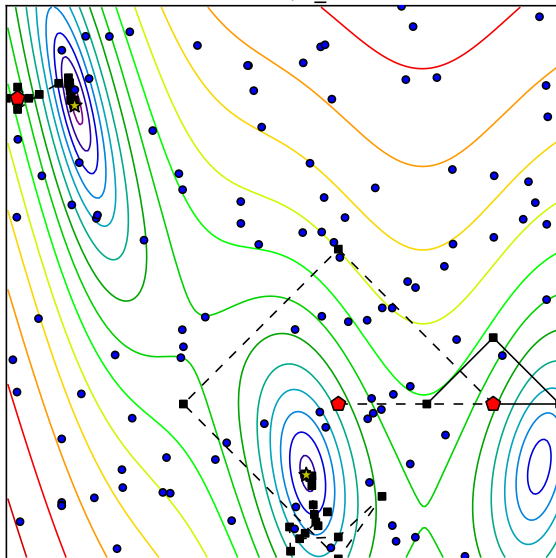
Iteration: 89; r_k : 0.263



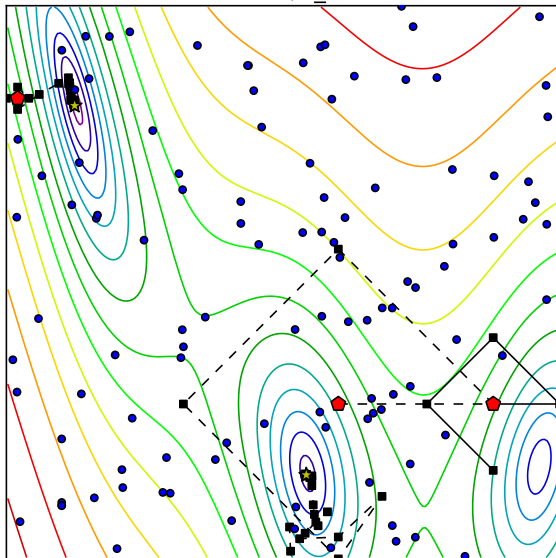
Iteration: 90; r_k : 0.262



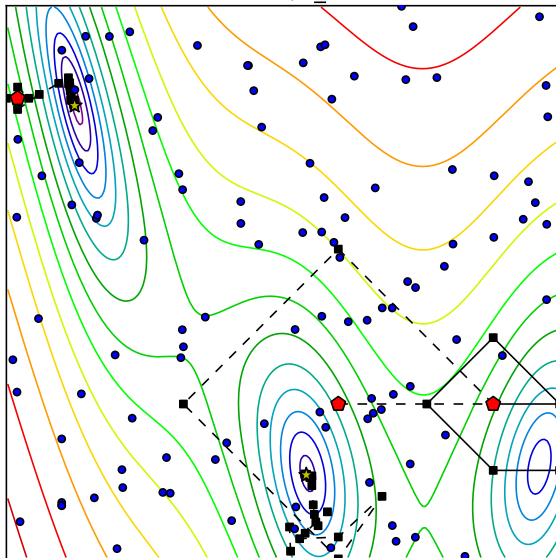
Iteration: 91; r_k : 0.261



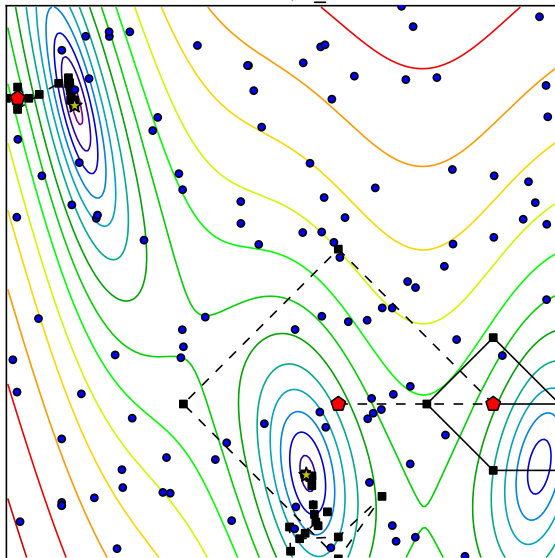
Iteration: 92; r_k : 0.260



Iteration: 93; r_k : 0.259

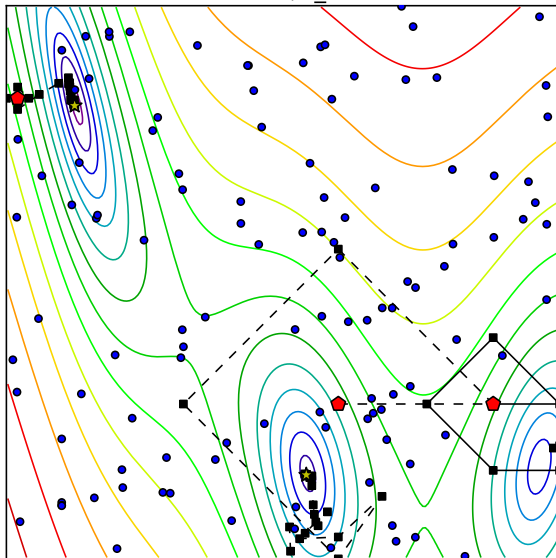


Iteration: 94; r_k : 0.258

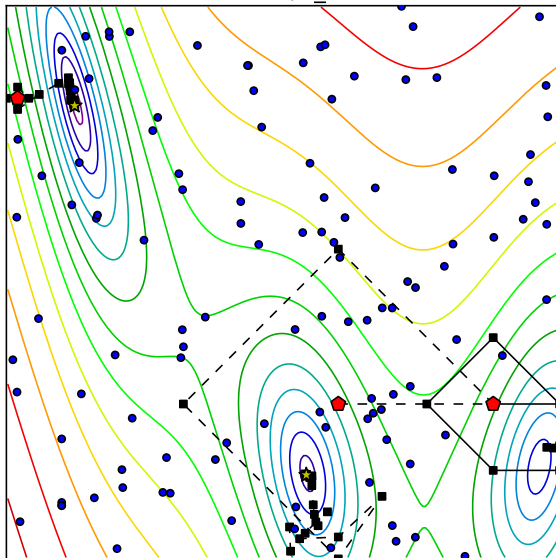


POSMM

Iteration: 95; r_k : 0.257

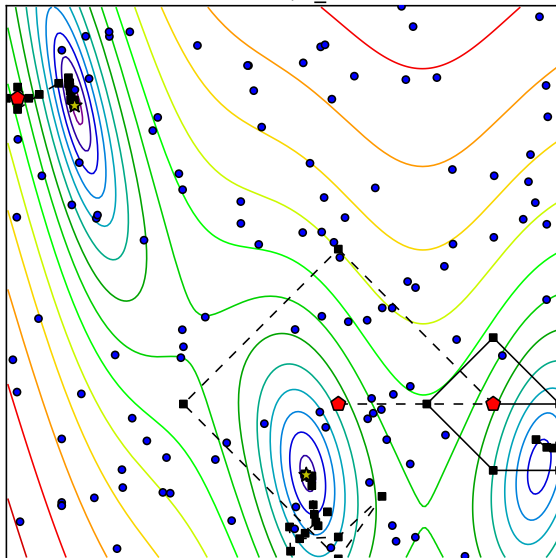


Iteration: 96; r_k : 0.256

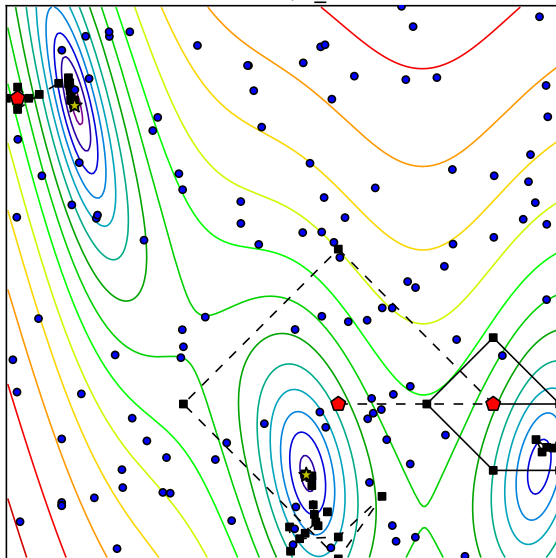


POSMM

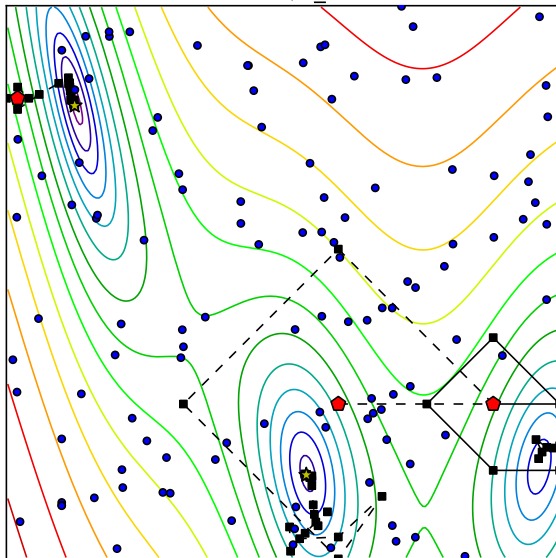
Iteration: 97; r_k : 0.255



Iteration: 98; r_k : 0.255



Iteration: 99; r_k : 0.254



Properties of the local optimization method

Necessary:

- ▶ Honors a starting point
- ▶ Honors bound constraints



Properties of the local optimization method

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ORBIT satisfies these [Wild, Regis, Shoemaker, SIAM-JOSC, 2008]

BOBYQA satisfies these [Powell, 2009]



Properties of the local optimization method

Necessary:

- ▶ Honors a starting point
- ▶ Honors bound constraints

ORBIT satisfies these [Wild, Regis, Shoemaker, SIAM-JOSC, 2008]

BOBYQA satisfies these [Powell, 2009]

Possibly beneficial:

- ▶ Can return multiple points of interest
- ▶ Reports solution quality/confidence at every iteration
- ▶ Can avoid certain regions in the domain
- ▶ Uses a history of past evaluations of f
- ▶ Uses additional points mid-run



Theorem

Given the same assumptions as MLSL, APOSSM will start a finite number of local optimization runs with probability 1.



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Assumption

There exists $K_0 < \infty$ so that for any K_0 consecutive iterations, there is a positive (bounded away from zero) probability of evaluating a point from the sample stream and each existing local optimization run.



Theorem

Given the same assumptions as MLSL, APOSSM will start a finite number of local optimization runs with probability 1.

Assumption

There exists $K_0 < \infty$ so that for any K_0 consecutive iterations, there is a positive (bounded away from zero) probability of evaluating a point from the sample stream and each existing local optimization run.

Theorem

Each $x^* \in X^*$ will almost surely be either identified in a finite number of evaluations or have a single local optimization run that is converging asymptotically to it.



Measuring Performance

GLODS Global & local optimization using direct search [Custódio, Madeira (JOGO, 2014)]

Direct Serial DIRECT [D. Finkel's MATLAB code]

pVTDirect Parallel DIRECT [He, Watson, Sosonkina (TOMS, 2009)]

Random Uniform sampling over domain (as a baseline)

POSMM

- ▶ Local optimization method

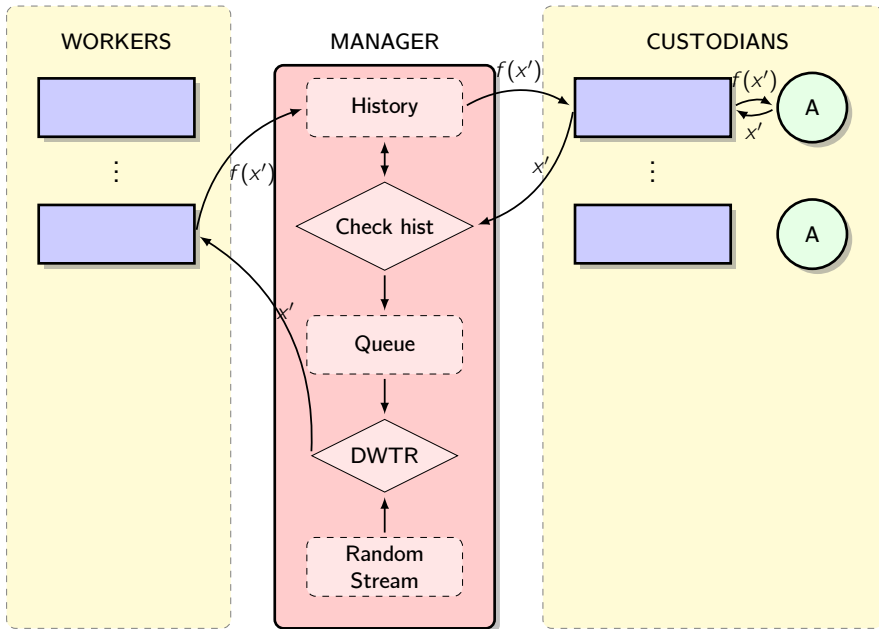
- ▶ ORBIT [Wild, Regis, & Shoemaker (SIAM JOSC, 2008)]

- ▶ BOBYQA [Powell, 2009]

- ▶ Initial sample size: $10n$

- ▶ Each method evaluates Direct's $2n + 1$ initial points.





(A) POSMM Diagram



(A) POSMM Manager

Algorithm 3: Manager Logic

```
while  $k < \text{fevalmax}$  do  
  MPI.recv(MPI.ANY_SOURCE)  
  if Received from Custodian then  
    Check  $H_k$   
    Add new point to  $Q_L$   
  if Received from Worker then  
    Update  $H_k$   
    Possibly get a Custodian working on the next point  
    Run decide_where_to_start  
    Possibly update  $Q_L$   
  if  $\text{sync} = \text{False}$  OR All Workers/Custodians are done then  
    Give from  $Q_L$  or  $\mathcal{R}_S$  to available worker(s)
```



POSMM

MLSL: (S2)–(S4)

$$\hat{x} \in S_k$$

- (S1) $\nexists x \in L_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (S2) $\nexists x \in S_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (S3) \hat{x} *has not started a local optimization run*
- (S4) \hat{x} *is at least μ from $\partial\mathcal{D}$ and ν from known local minima*

POSMM: (S1)–(S4), (L1)–(L6)

$$\hat{x} \in L_k$$

- (L1) $\nexists x \in L_k$
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (L2) $\nexists x \in S_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (L3) \hat{x} has not started a local optimization run
- (L4) \hat{x} is at least μ from $\partial\mathcal{D}$ and ν from known local minima
- (L5) \hat{x} is not in an active local optimization run and has not been ruled stationary
- (L6) $\exists r_k$ -descent path in H_k from some $x \in S_k$ satisfying (S2-S4) to \hat{x}



Measuring Performance

Notation:

Let X^* be the set of all local minima of f .

Let $f_{(i)}^*$ be the i th smallest value $\{f(x^*) | x^* \in X^*\}$.

Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The global minimum has been found at a level $\tau > 0$ at batch k if an algorithm it has found a point \hat{x} satisfying:

$$f(\hat{x}) - f_{(1)}^* \leq (1 - \tau) \left(f(x_0) - f_{(1)}^* \right),$$

where x_0 is the starting point for problem p .



Measuring Performance

Notation:

Let X^* be the set of all local minima of f .

Let $f_{(i)}^*$ be the i th smallest value $\{f(x^*) | x^* \in X^*\}$.

Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The j best local minima have been found at a level $\tau > 0$ at batch k if:

$$\left| \left\{ x_{(1)}^*, \dots, x_{(\underline{j}-1)}^* \right\} \cap \left\{ x_{(i)}^* : \exists x \in H_k \text{ with } \|x - x_{(i)}^*\| \leq r_n(\tau) \right\} \right| = \underline{j} - 1$$

&

$$\left| \left\{ x_{(\underline{j})}^*, \dots, x_{(\bar{j})}^* \right\} \cap \left\{ x_{(i)}^* : \exists x \in H_k \text{ with } \|x - x_{(i)}^*\| \leq r_n(\tau) \right\} \right| \geq \bar{j} - \underline{j} + 1,$$

where \underline{j} and \bar{j} are the smallest and largest integers such that

$$f_{(\underline{j})}^* = f_{(\underline{j})}^* = f_{(\underline{j})}^* \text{ and where } r_n(\tau) = \sqrt[n]{\frac{\tau \text{vol}(\mathcal{D}) \Gamma(\frac{n}{2} + 1)}{\pi^{n/2}}}.$$



Problems considered

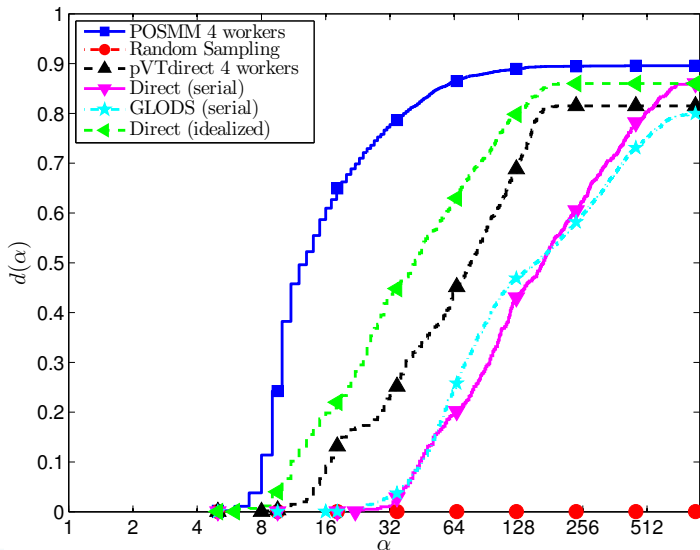
GKLS problem generator [Gaviano et al., “Algorithm 829” (TOMS, 2003)]

- ▶ 600 synthetic problems with known local minima
- ▶ $n = 2, \dots, 7$
- ▶ 10 local minima in the unit cube with a unique global minimum
- ▶ 100 problems for each dimension
- ▶ 5 replications (different seeds) for each problem
- ▶ 5000 evaluations



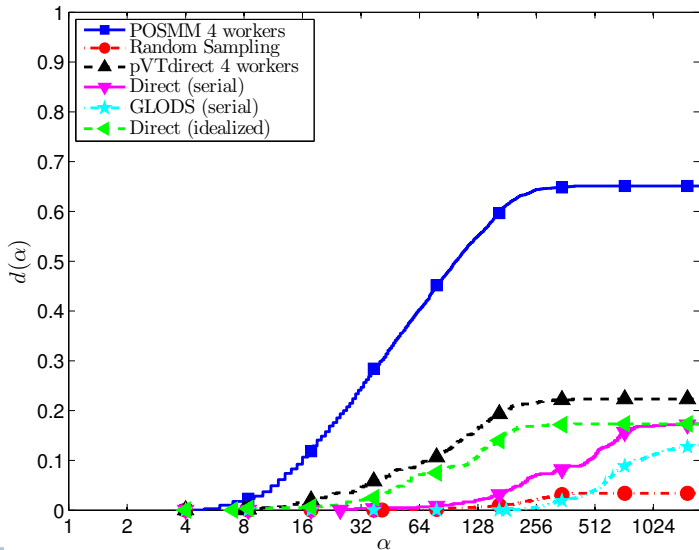
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



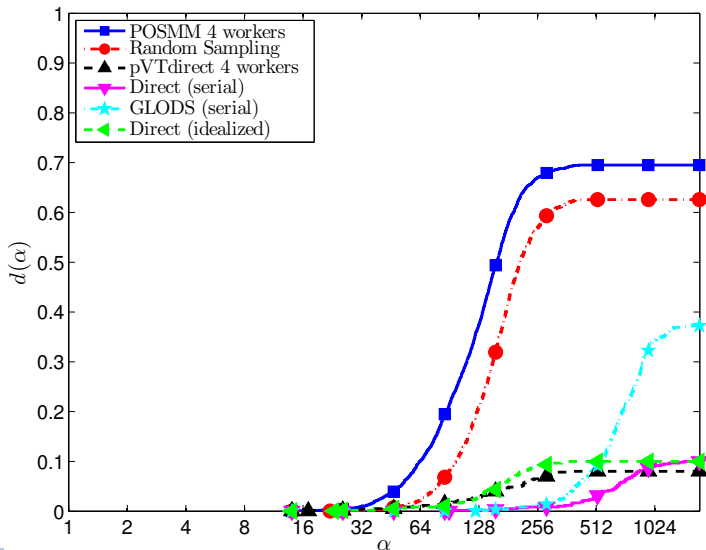
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



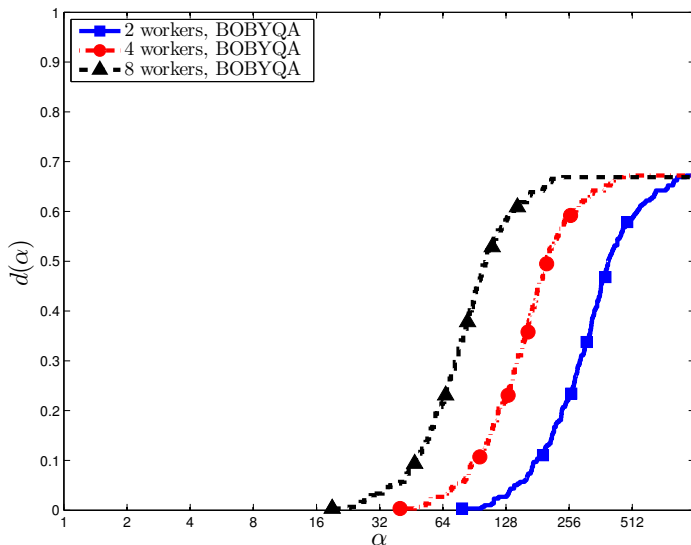
Data Profiles

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



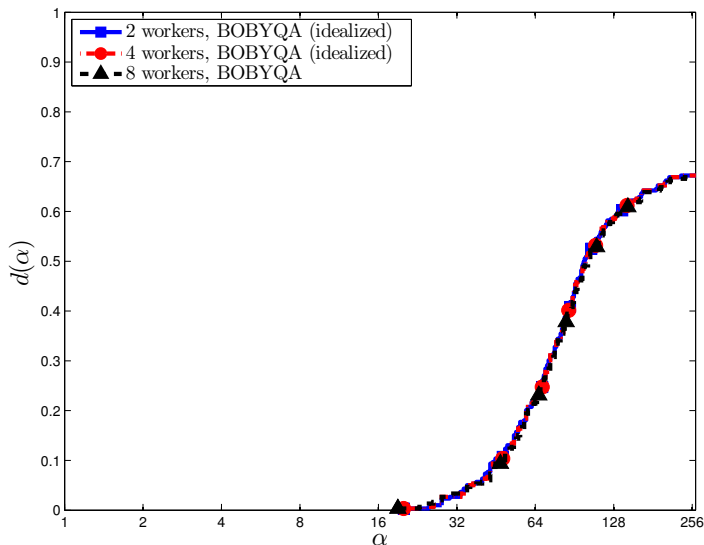
Data Profiles (increasing workers)

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



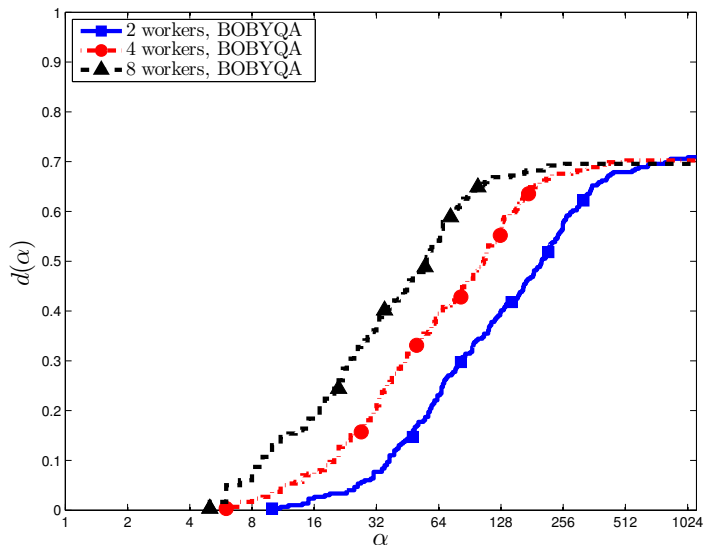
Data Profiles (increasing workers)

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



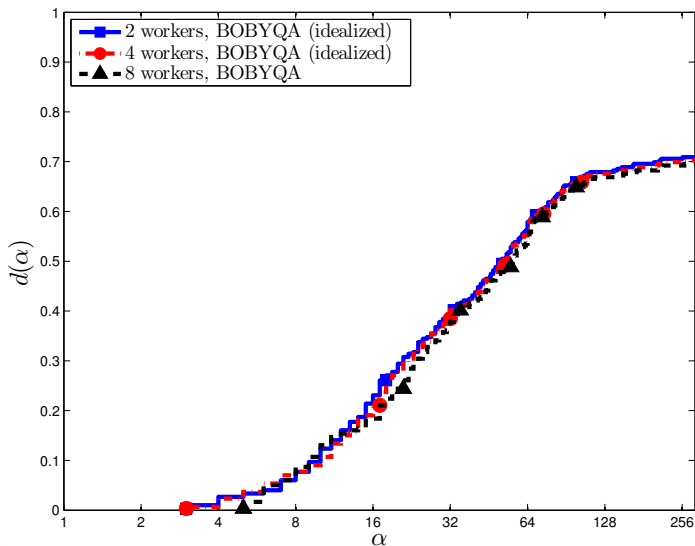
Data Profiles (increasing workers)

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



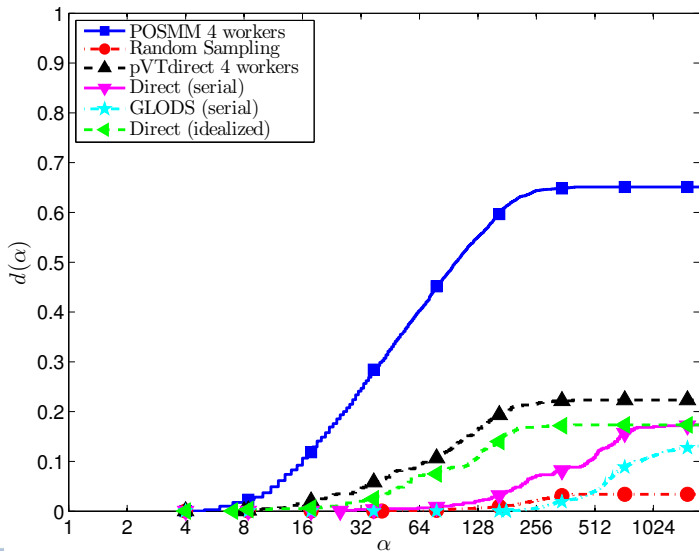
Data Profiles (increasing workers)

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



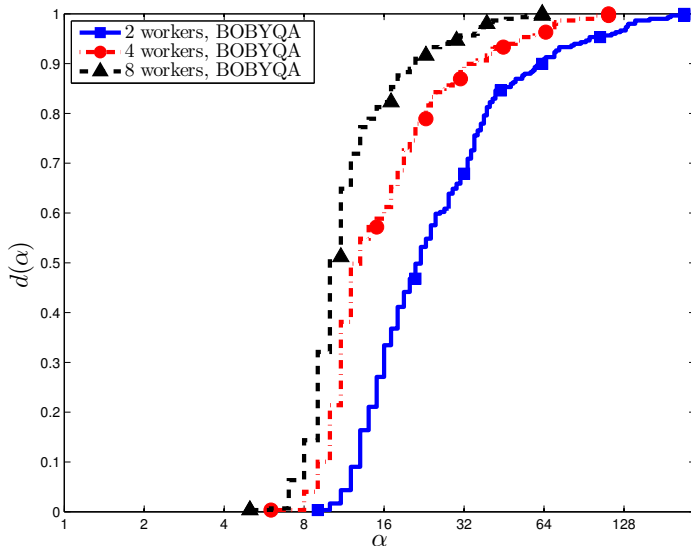
Data Profiles (increasing workers)

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



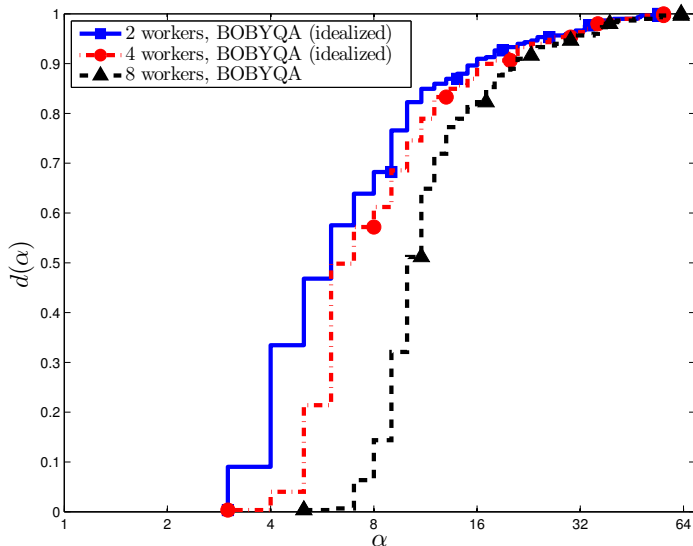
Data Profiles (increasing workers)

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



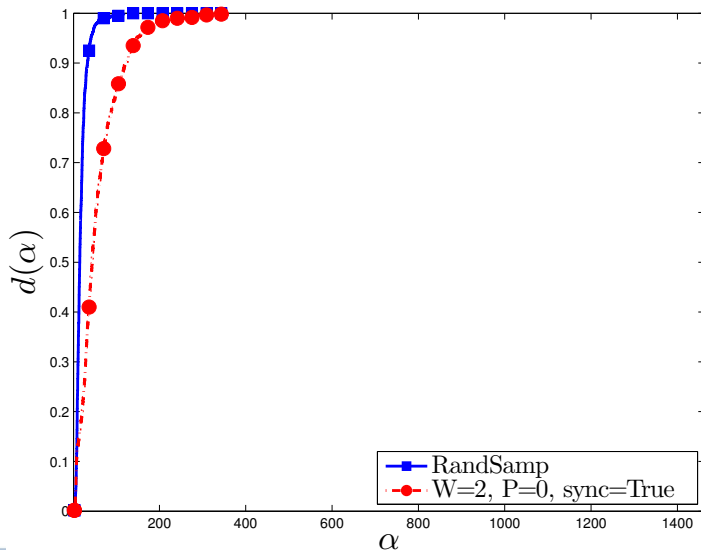
Data Profiles (increasing workers)

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



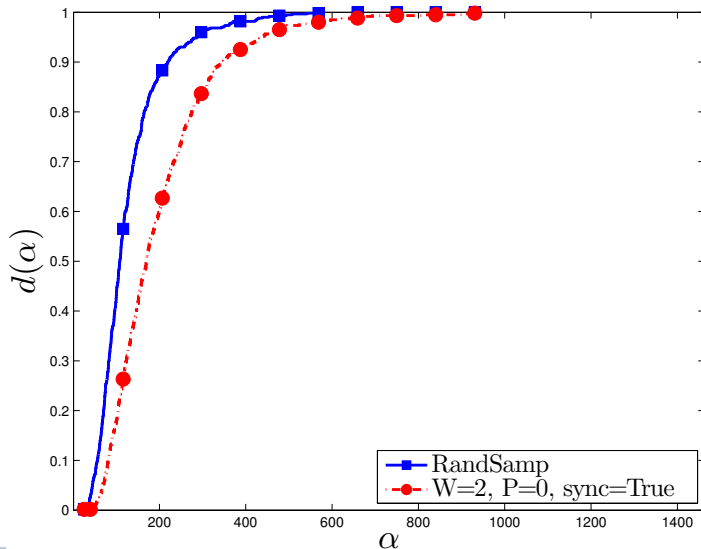
Data Profiles (vs random)

Within $\sqrt[n]{\frac{10^{-1}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 10 best minima



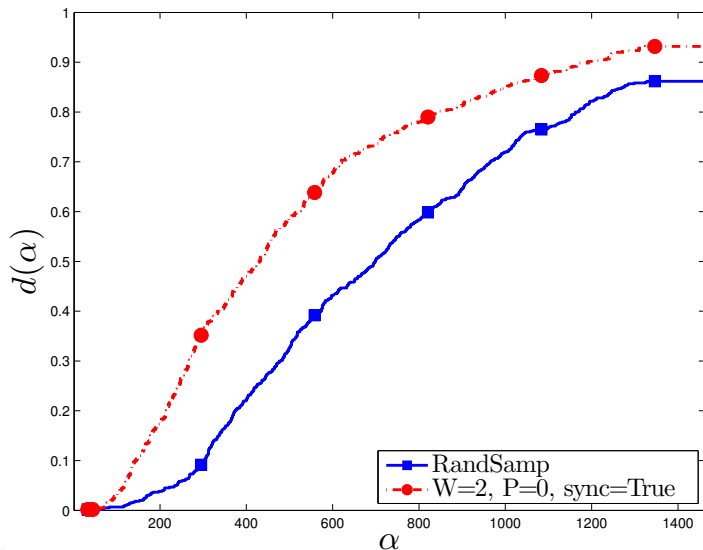
Data Profiles (vs random)

Within $\sqrt[n]{\frac{10^{-2}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 10 best minima



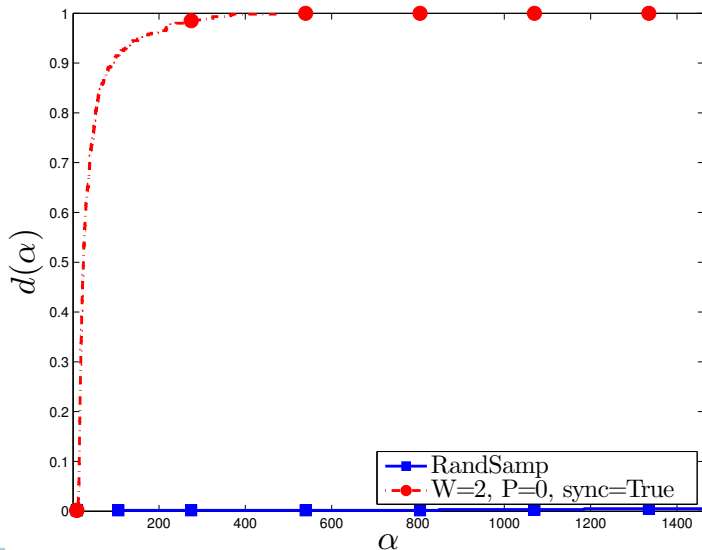
Data Profiles (vs random)

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima

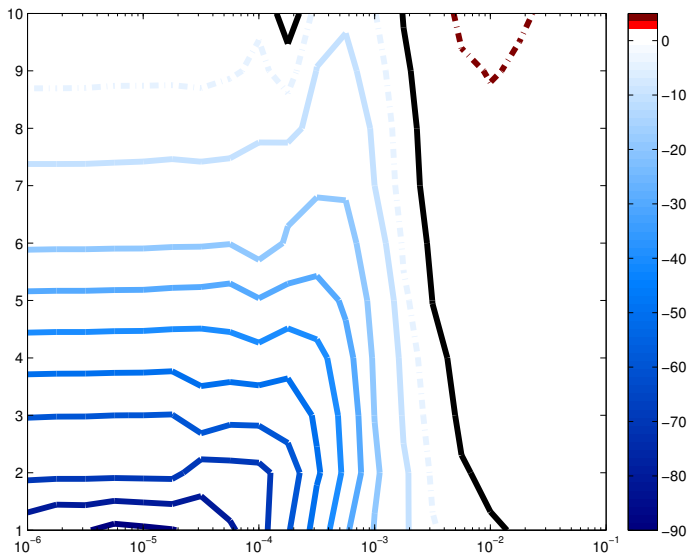


Data Profiles (vs random)

Within $\sqrt[n]{\frac{10^{-6}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of best minimum

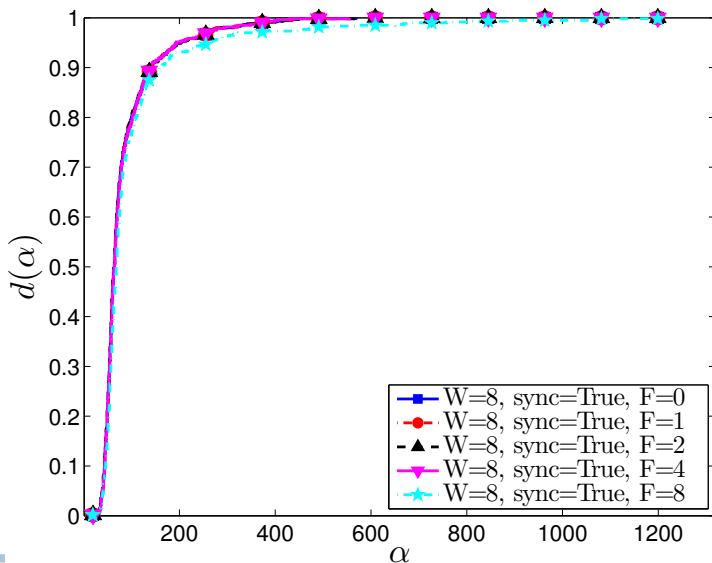


Data Profiles (vs random)



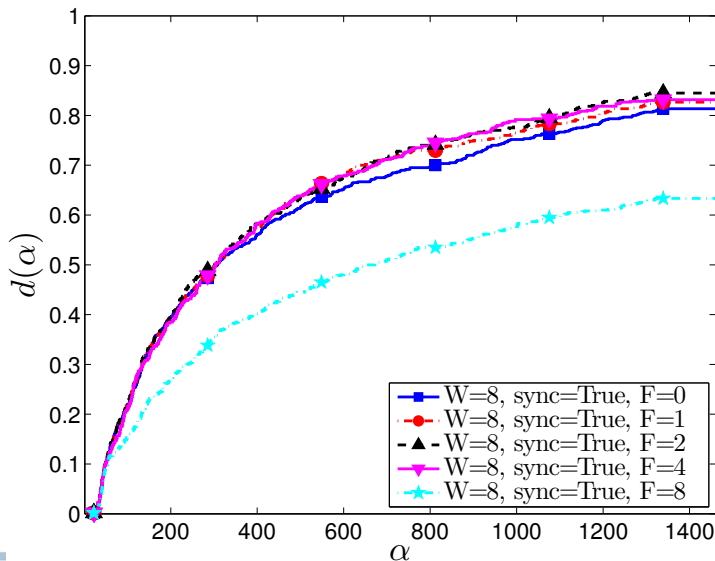
Data Profiles (forcing runs)

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



Data Profiles (forcing runs)

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



Closing Remarks

- ▶ Concurrent function evaluations can locate multiple minima while efficiently finding a global minimum.



Closing Remarks

- ▶ Concurrent function evaluations can locate multiple minima while efficiently finding a global minimum.
- ▶ The ability to find many minima scales well with the number of workers.

Questions:

- ▶ Finding (or designing) the best local solver for our framework?
- ▶ Best way to process the queue?



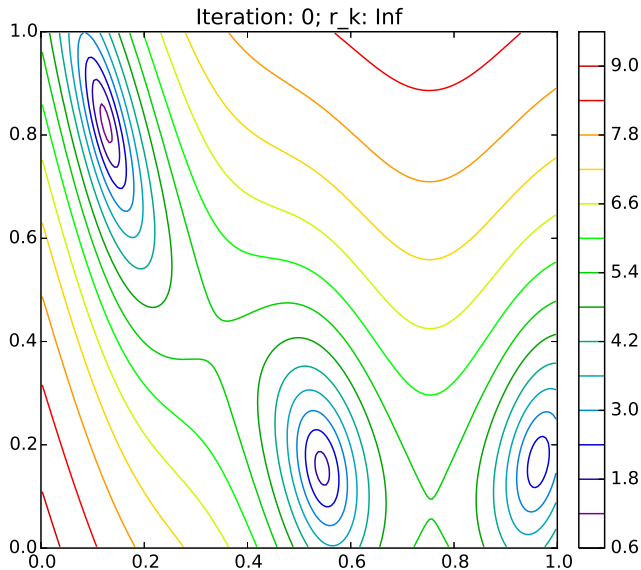
(A)POSMM Manager

Algorithm 3: Manager Logic

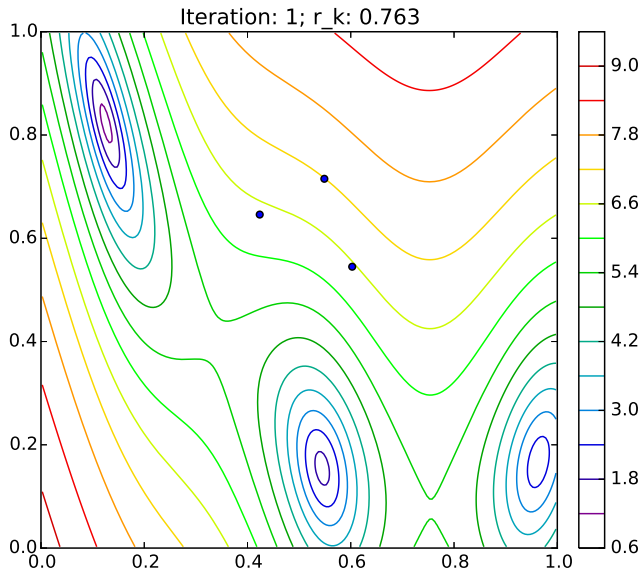
```
while  $k < \text{fevalmax}$  do  
  MPI.recv(MPI.ANY_SOURCE)  
  if Received from Custodian then  
    Check  $H_k$   
    Add new point to  $Q_L$   
  if Received from Worker then  
    Update  $H_k$   
    Possibly get a Custodian working on the next point  
    Run decide_where_to_start  
    Possibly update  $Q_L$   
  if sync=False OR All Workers/Custodians are done then  
    Give from  $Q_L$  or  $\mathcal{R}_S$  to available worker(s)
```



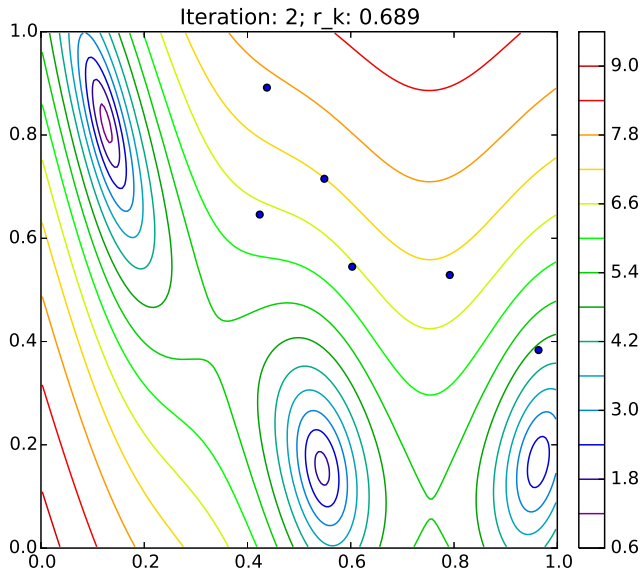
Pausing runs



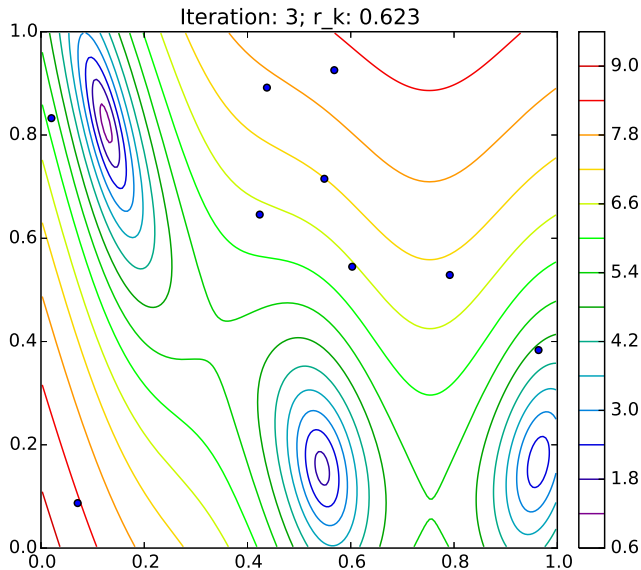
Pausing runs



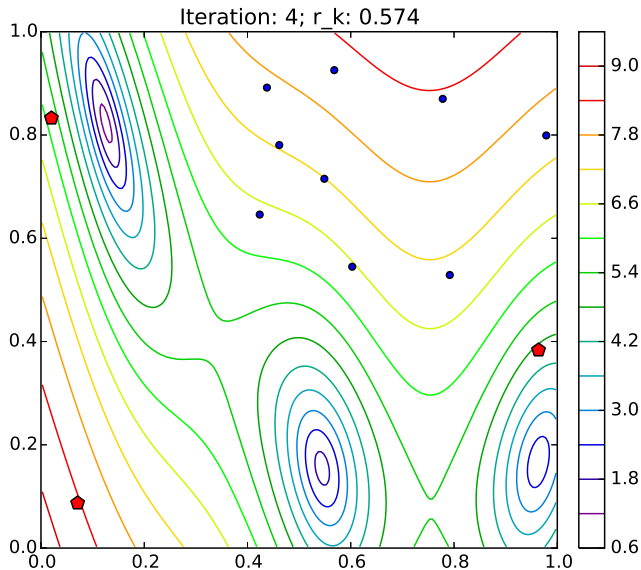
Pausing runs



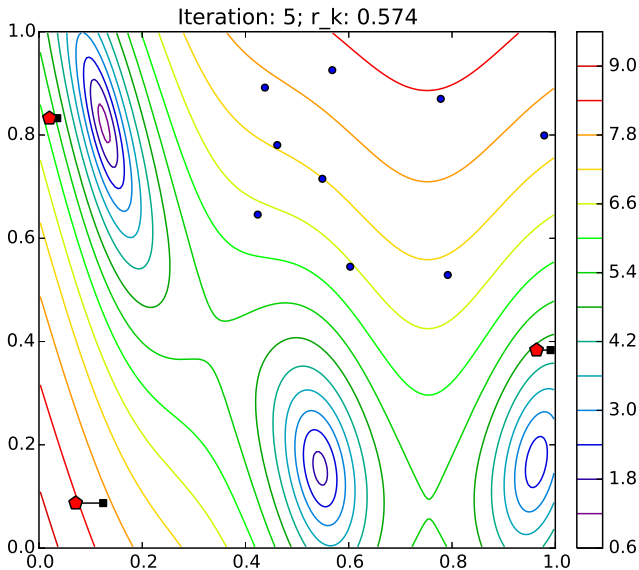
Pausing runs



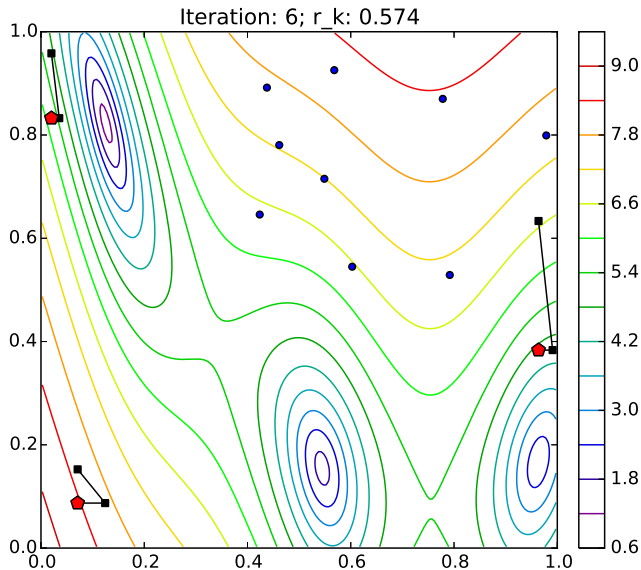
Pausing runs



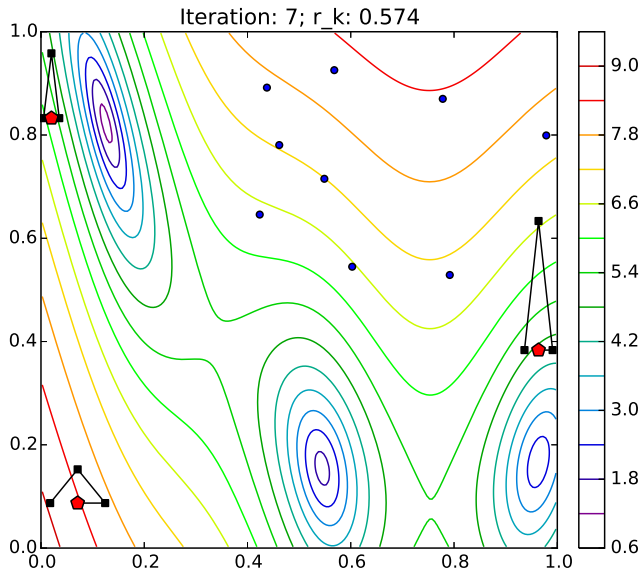
Pausing runs



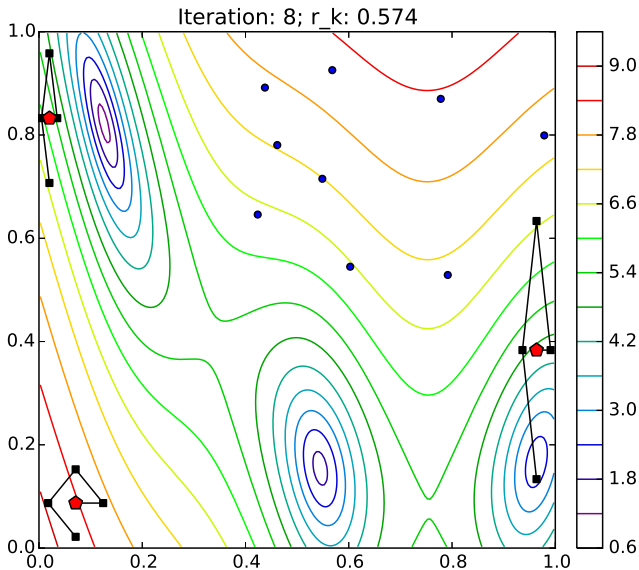
Pausing runs



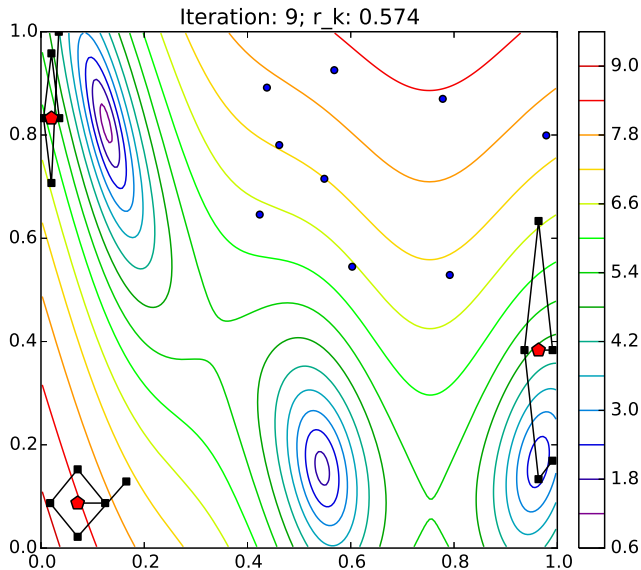
Pausing runs



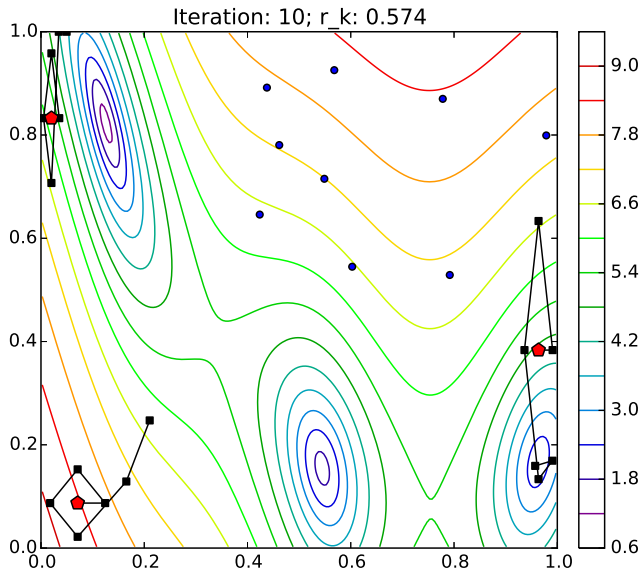
Pausing runs



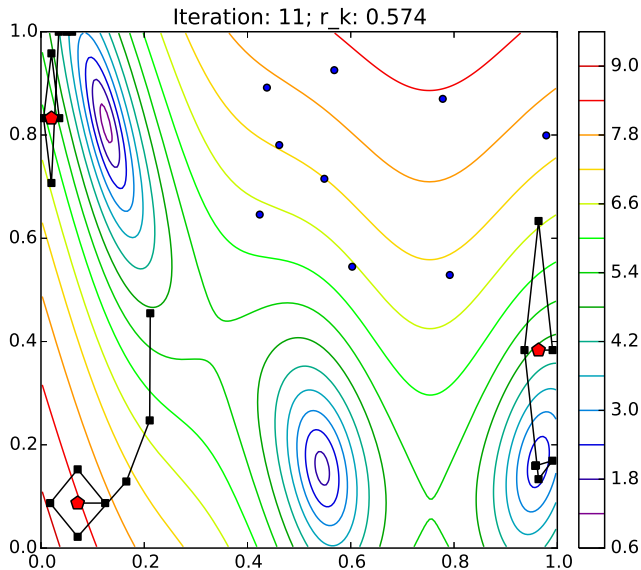
Pausing runs



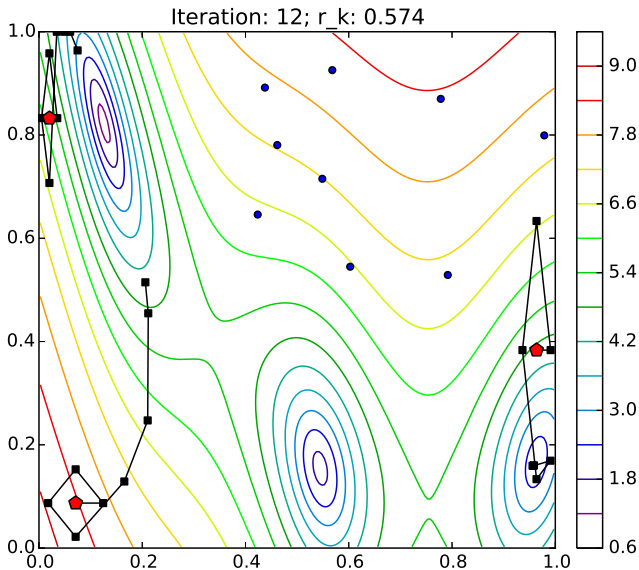
Pausing runs



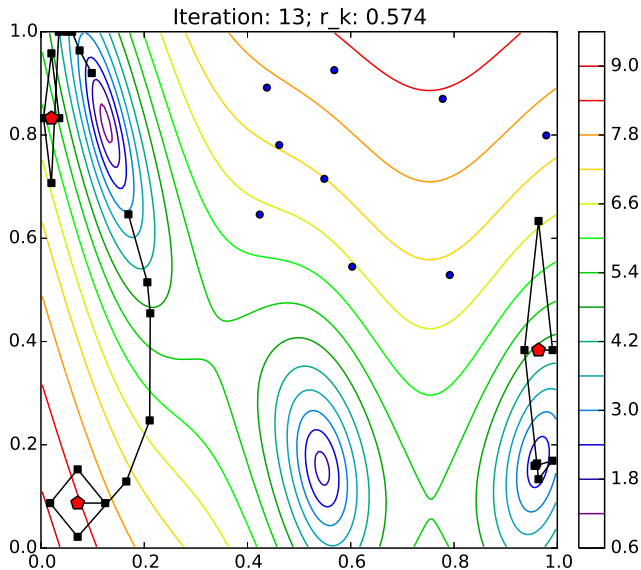
Pausing runs



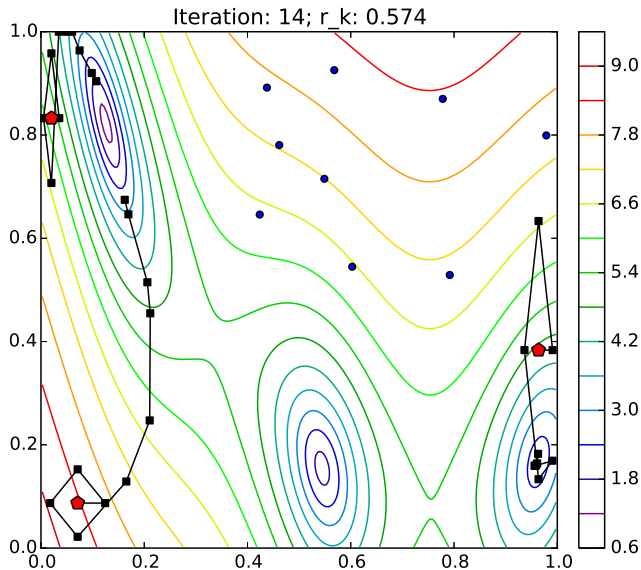
Pausing runs



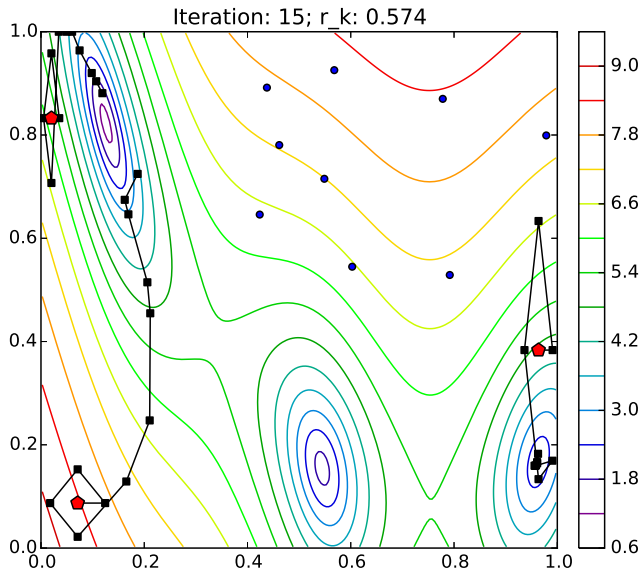
Pausing runs



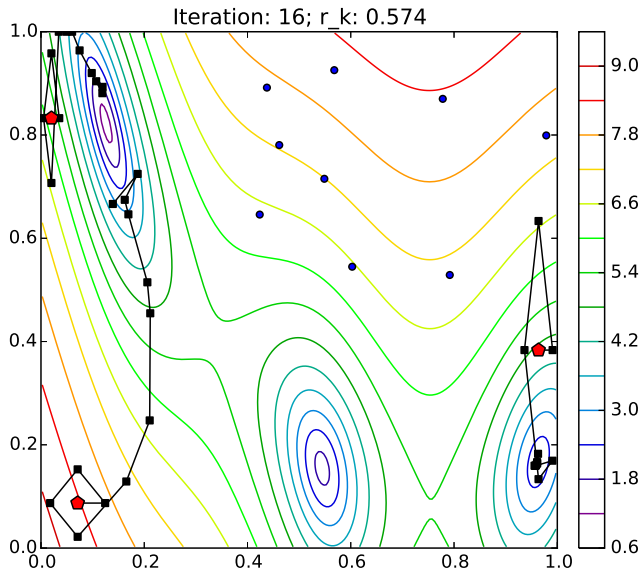
Pausing runs



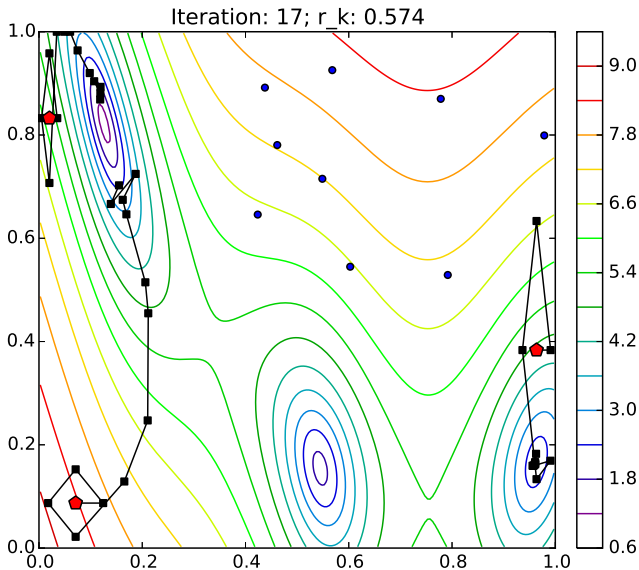
Pausing runs



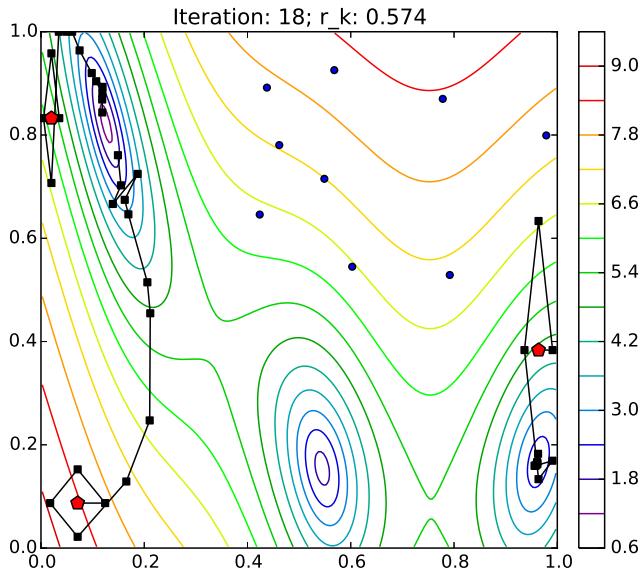
Pausing runs



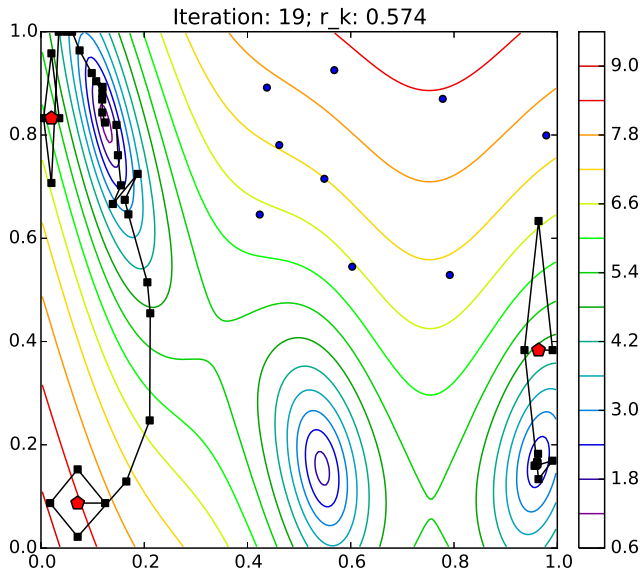
Pausing runs



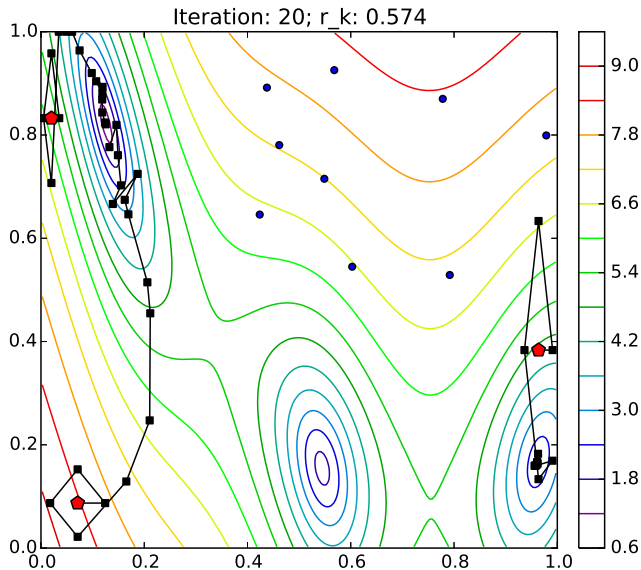
Pausing runs



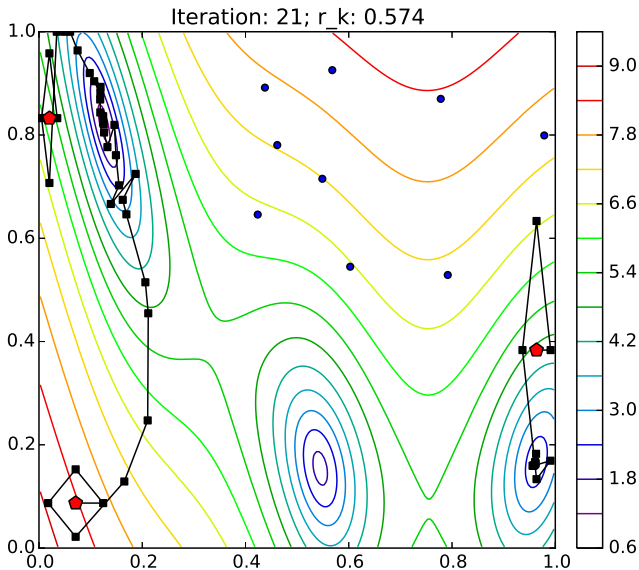
Pausing runs



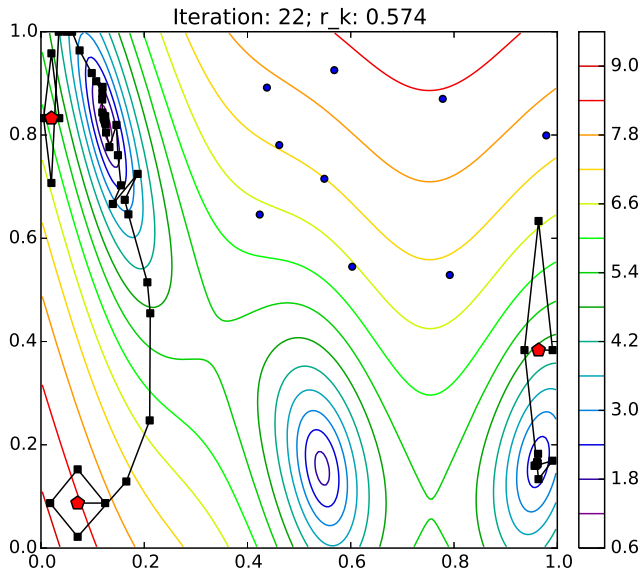
Pausing runs



Pausing runs

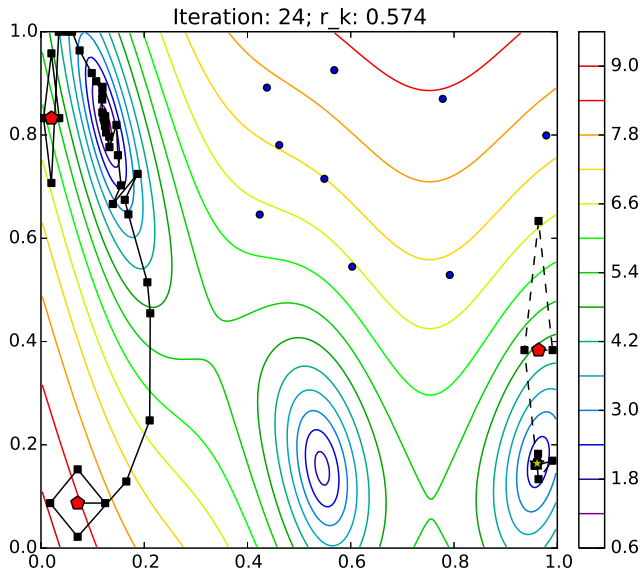


Pausing runs

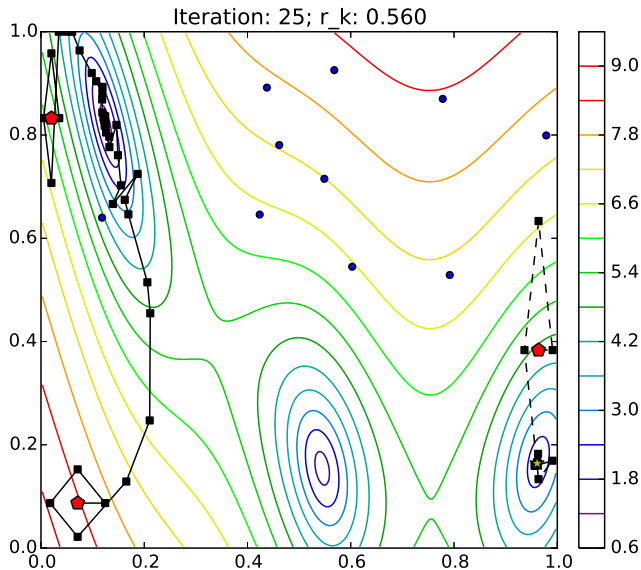




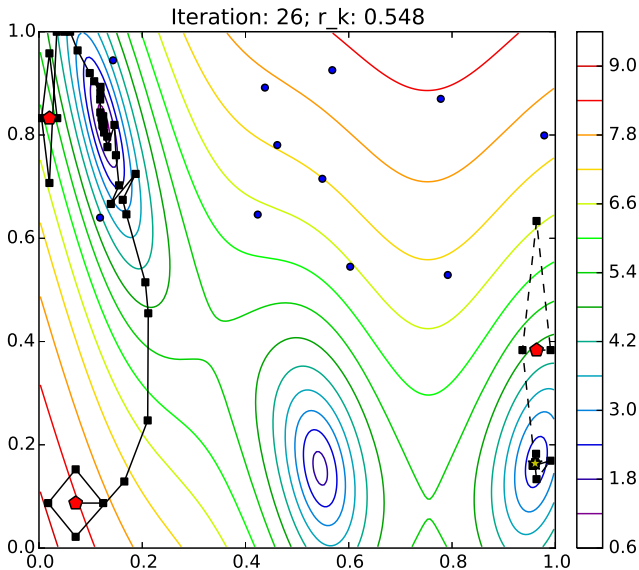
Pausing runs



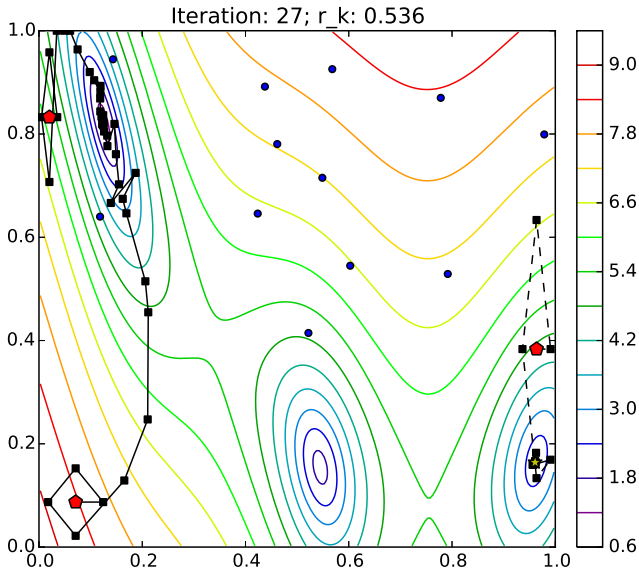
Pausing runs



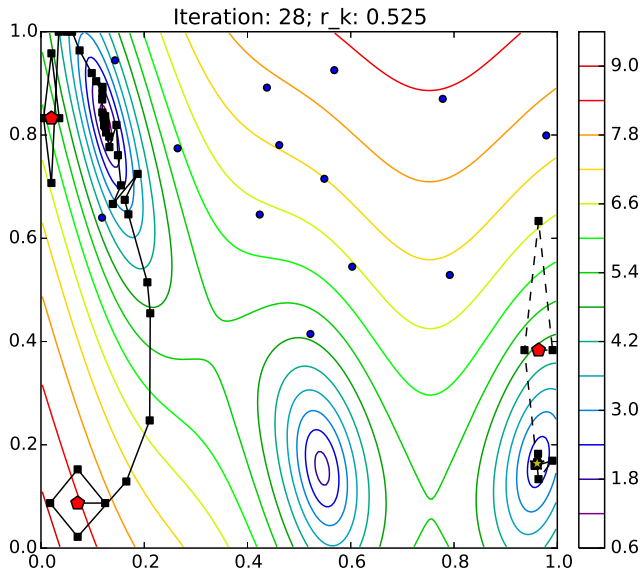
Pausing runs



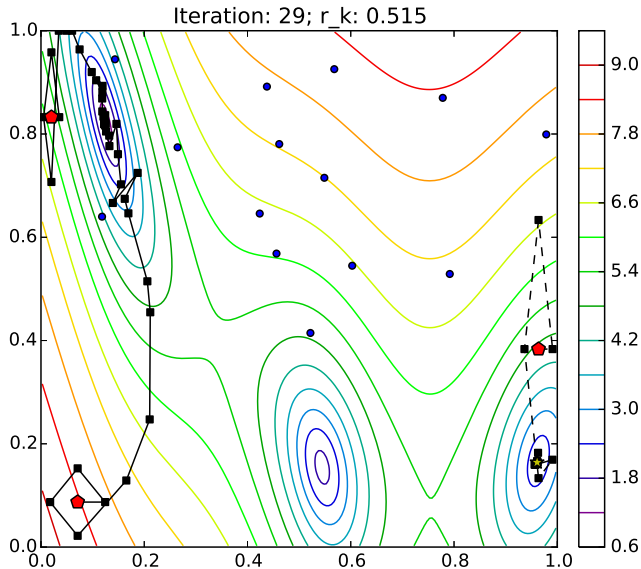
Pausing runs



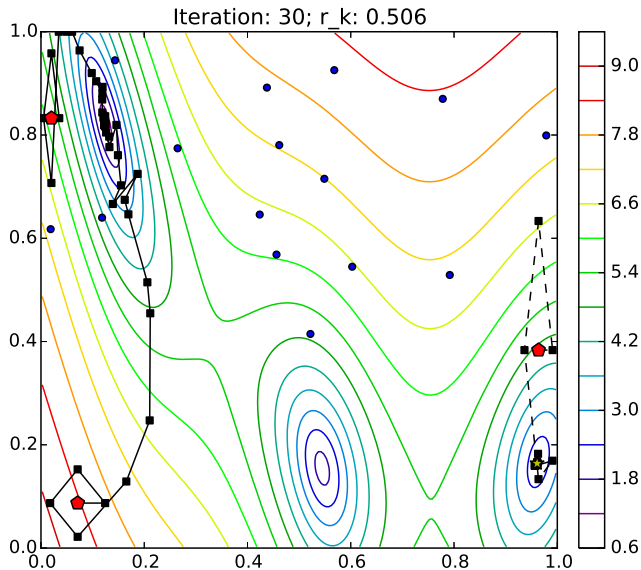
Pausing runs



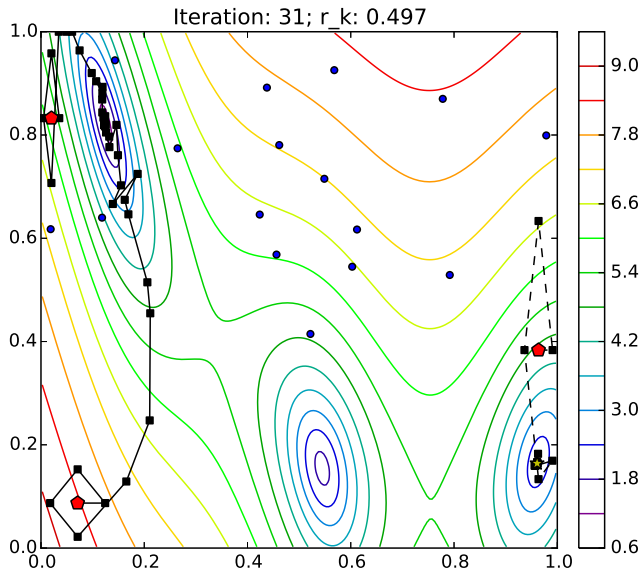
Pausing runs



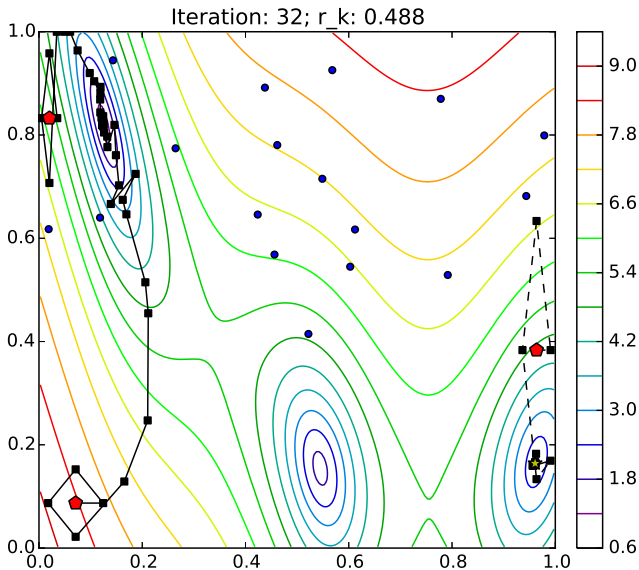
Pausing runs



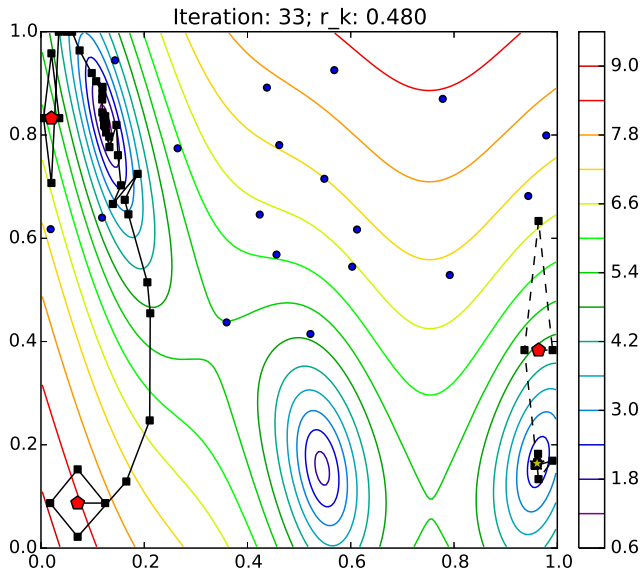
Pausing runs



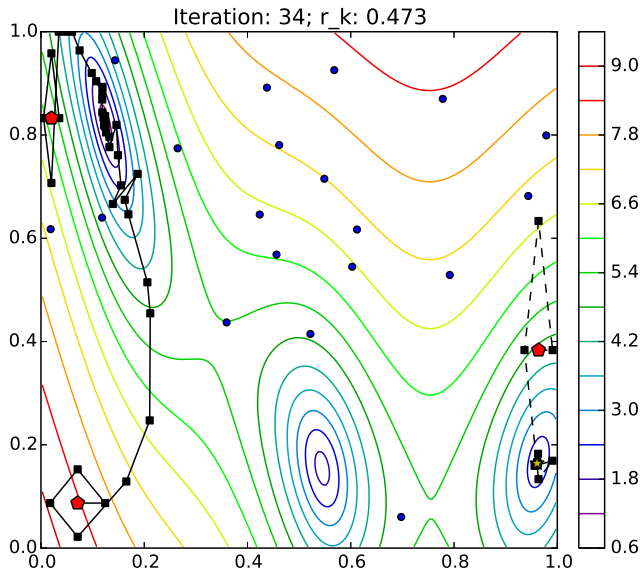
Pausing runs



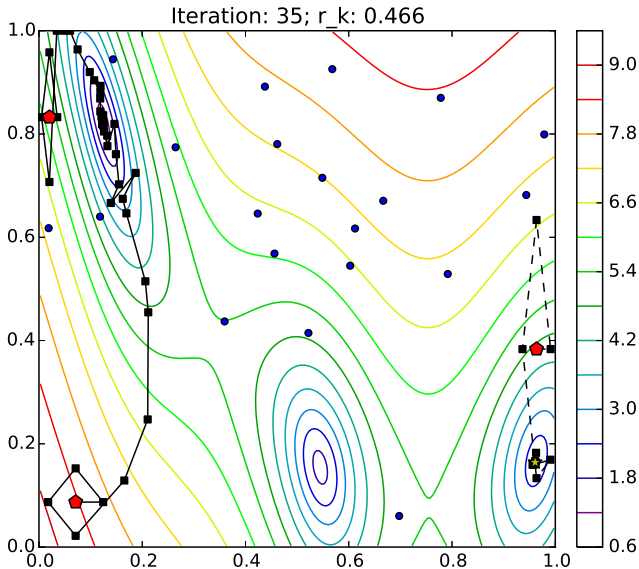
Pausing runs



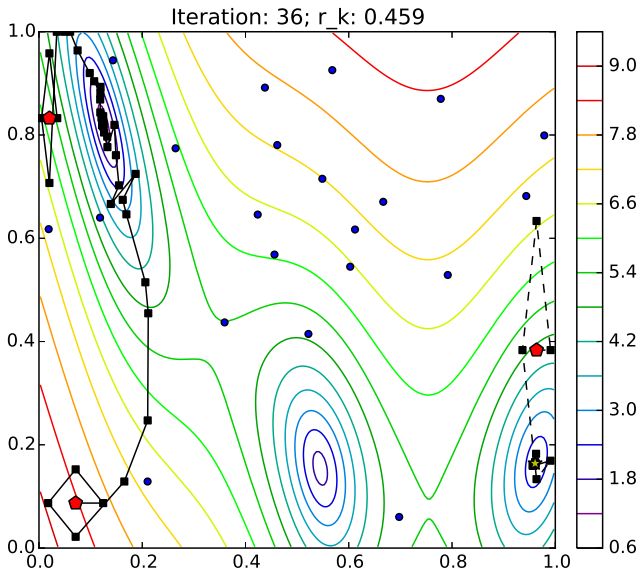
Pausing runs



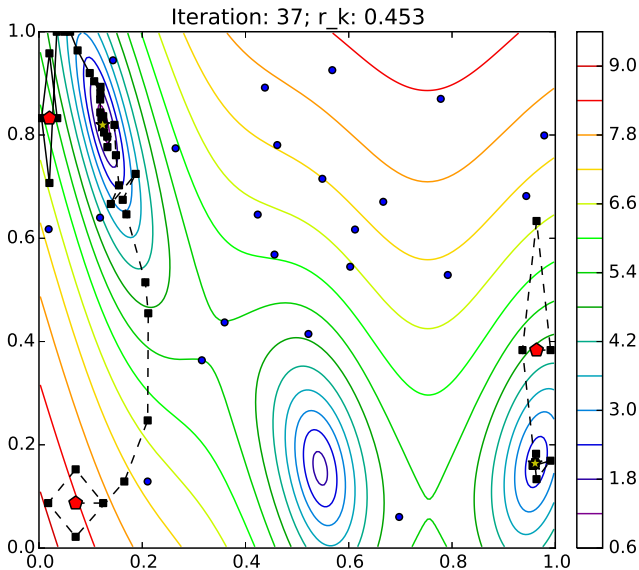
Pausing runs



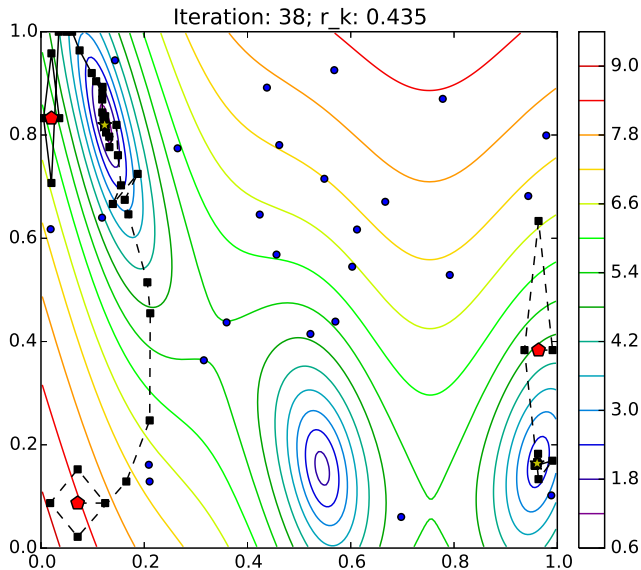
Pausing runs



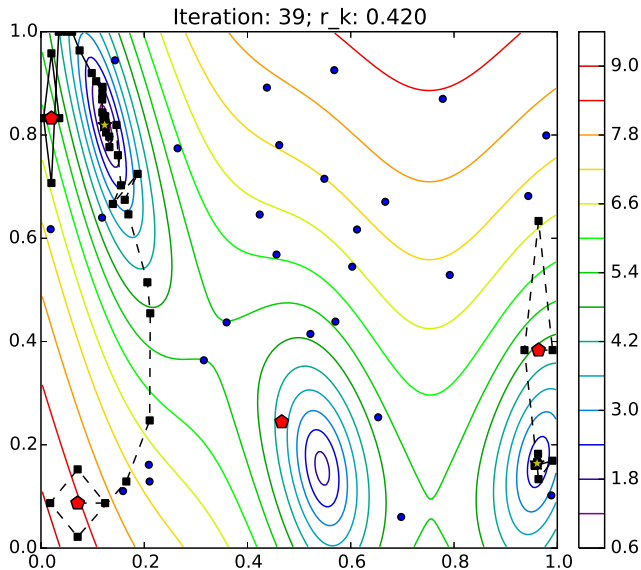
Pausing runs



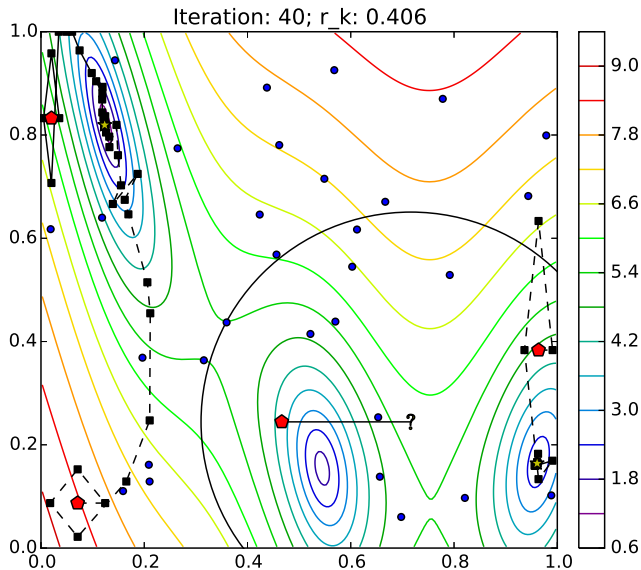
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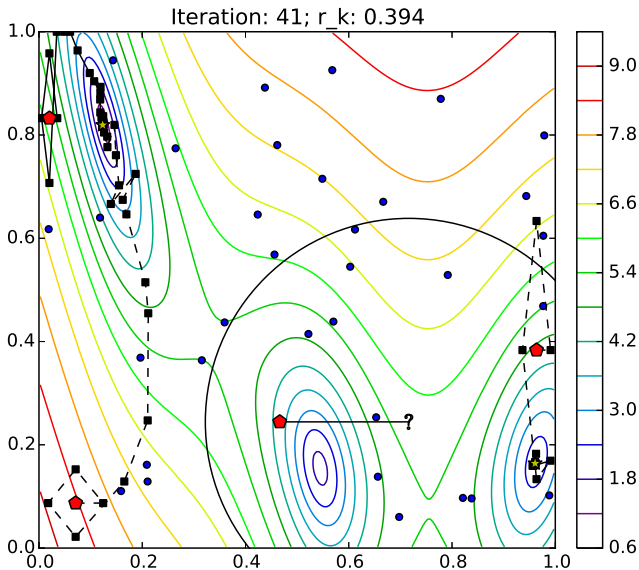
Pausing runs



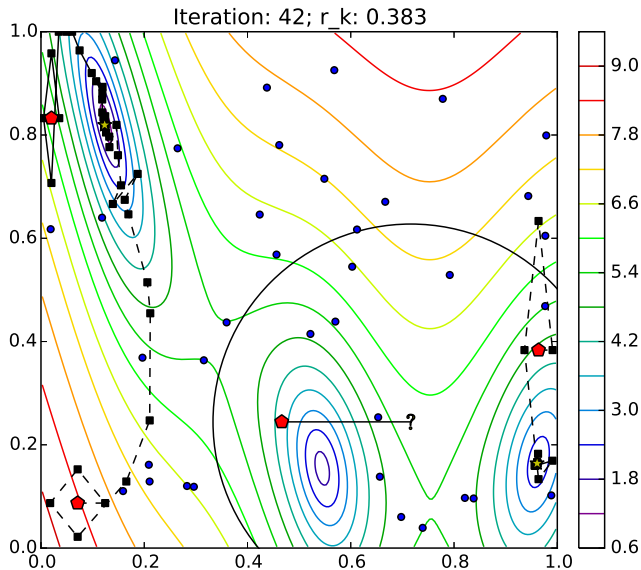
Pausing runs



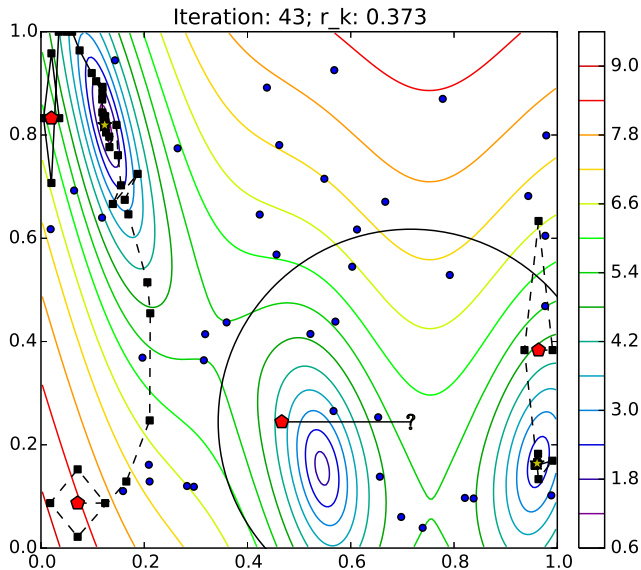
Pausing runs



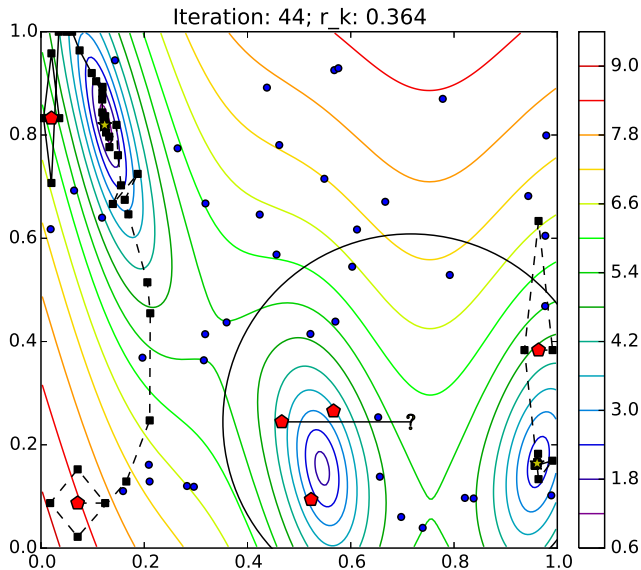
Pausing runs



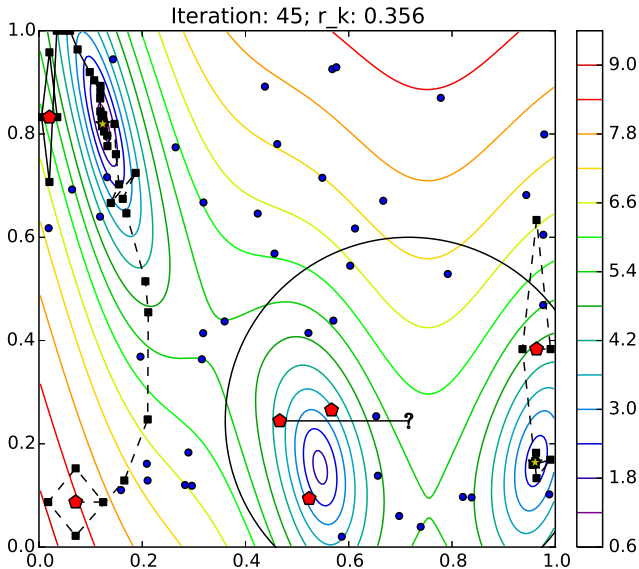
Pausing runs



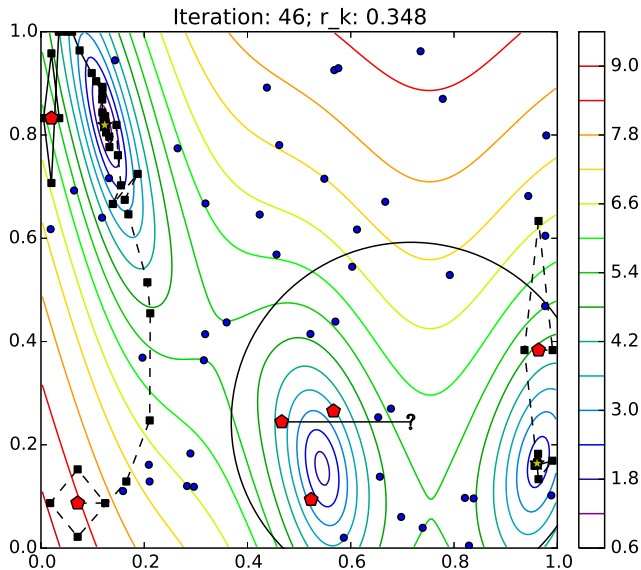
Pausing runs



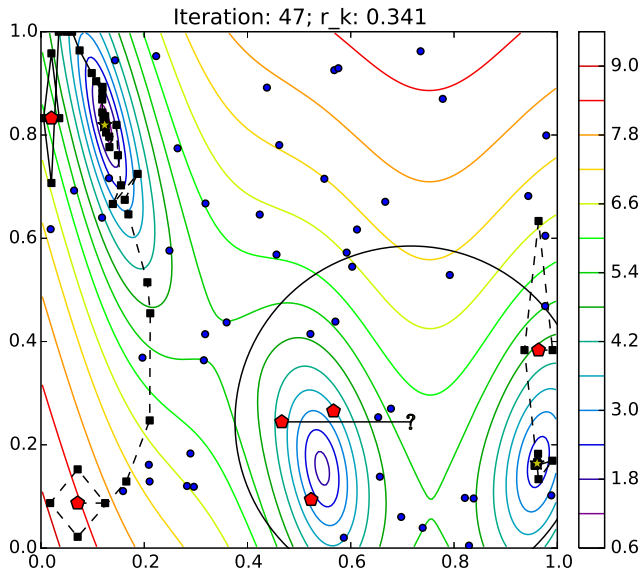
Pausing runs



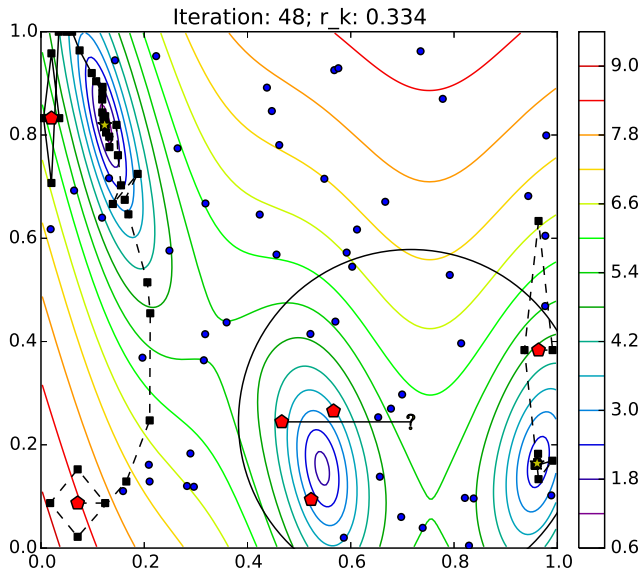
Pausing runs



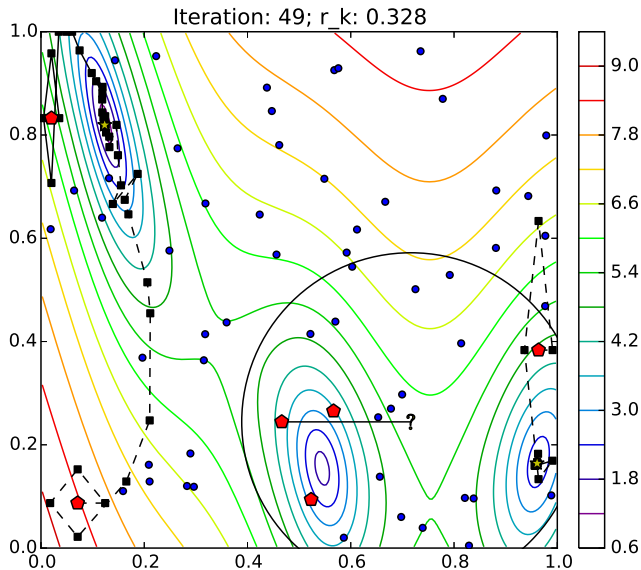
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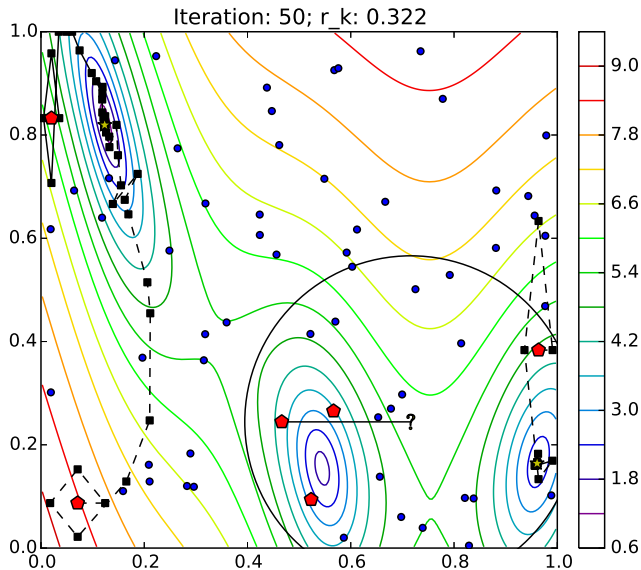
Pausing runs



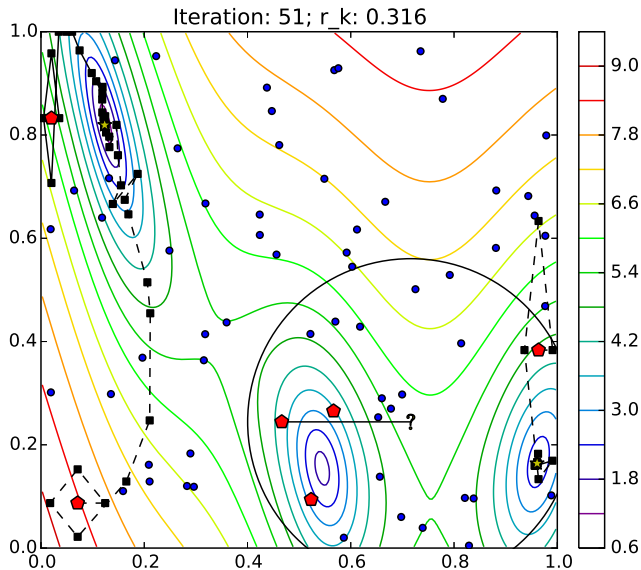
Pausing runs



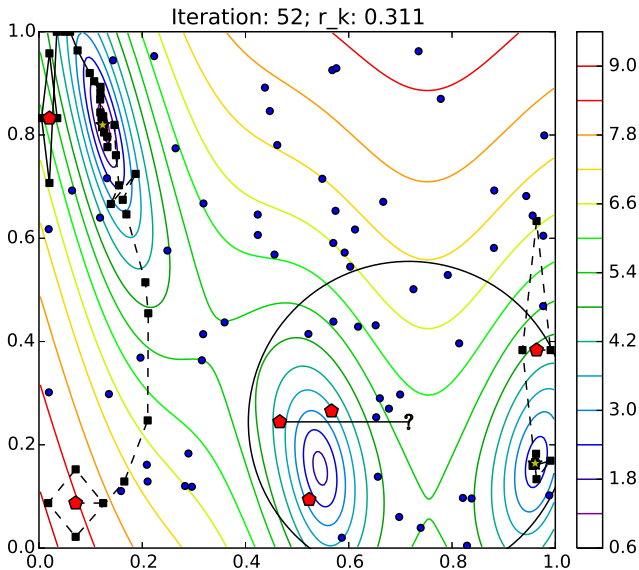
Pausing runs



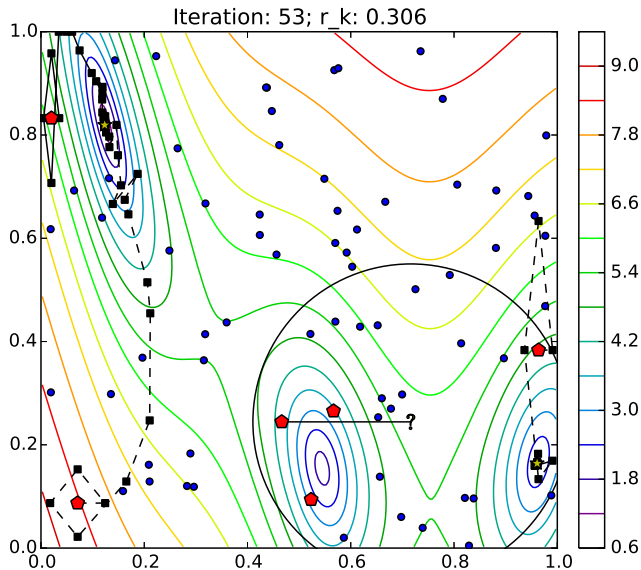
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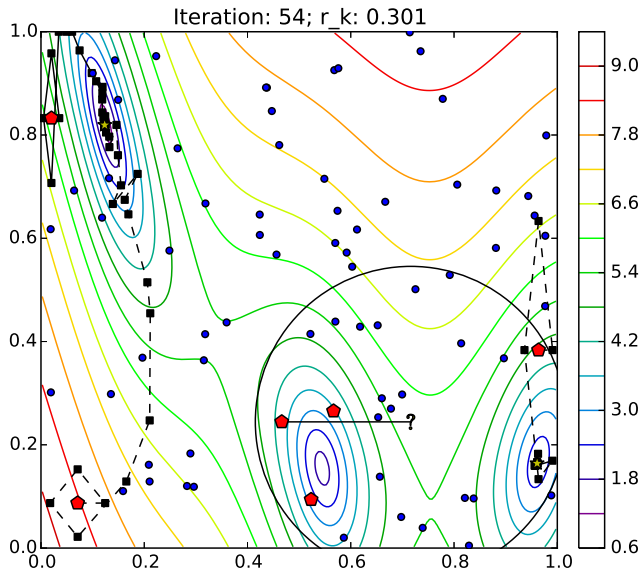
Pausing runs



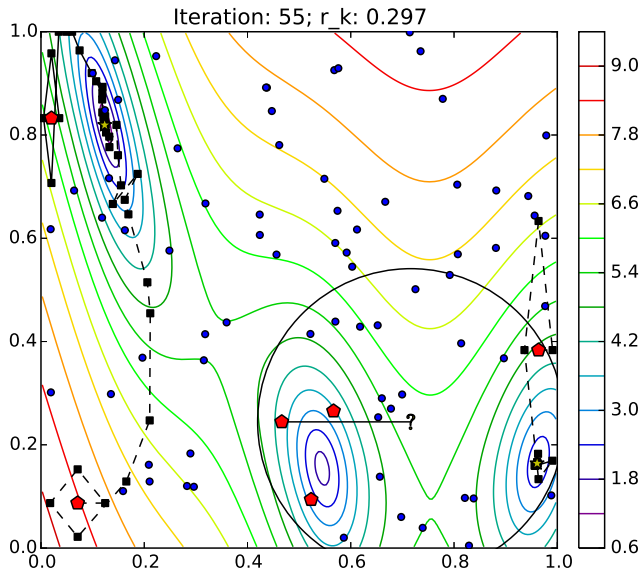
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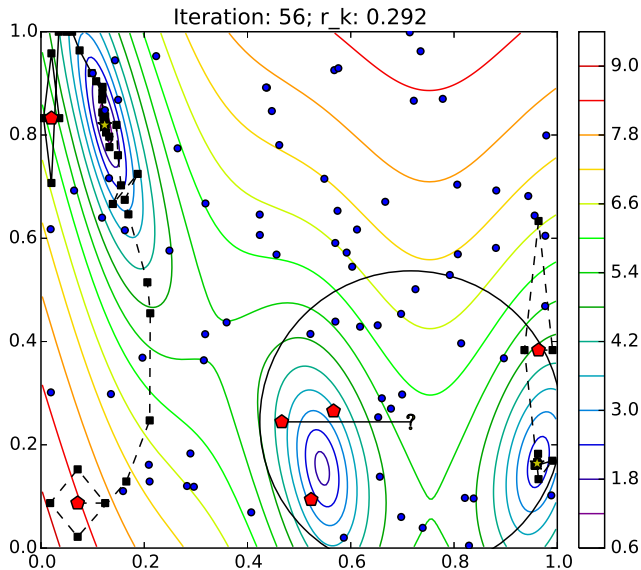
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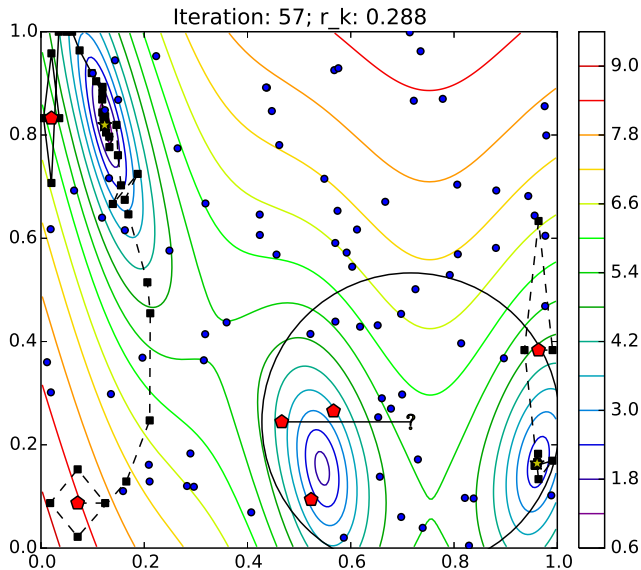
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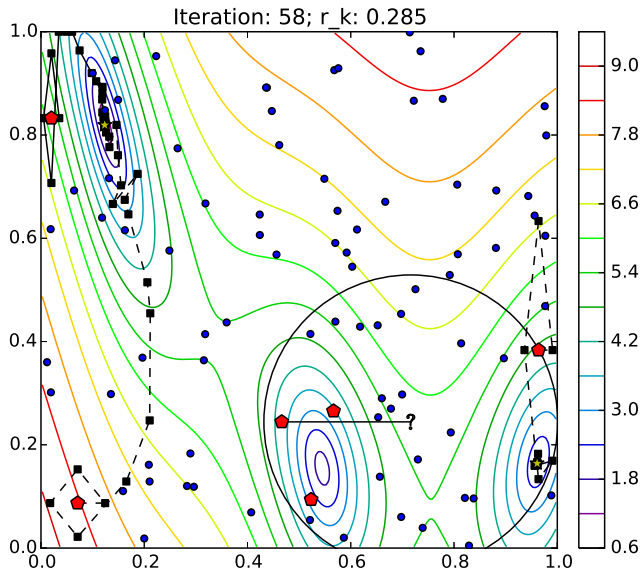
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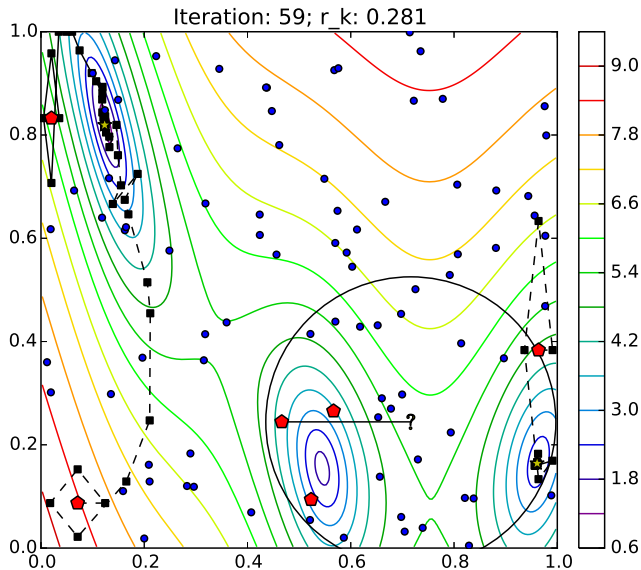
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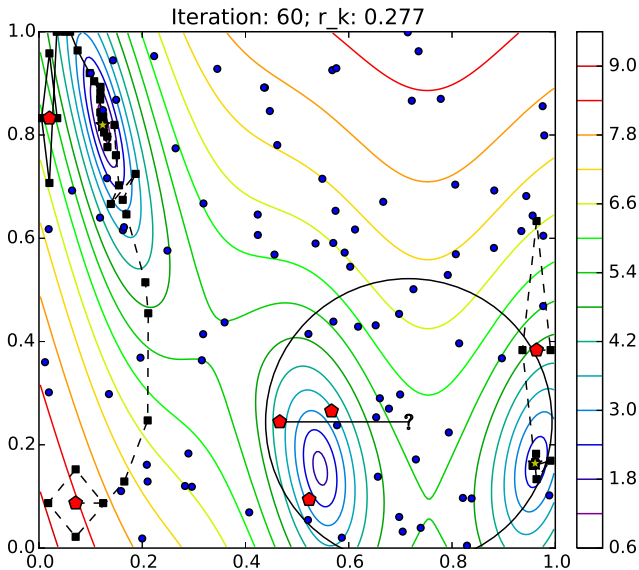
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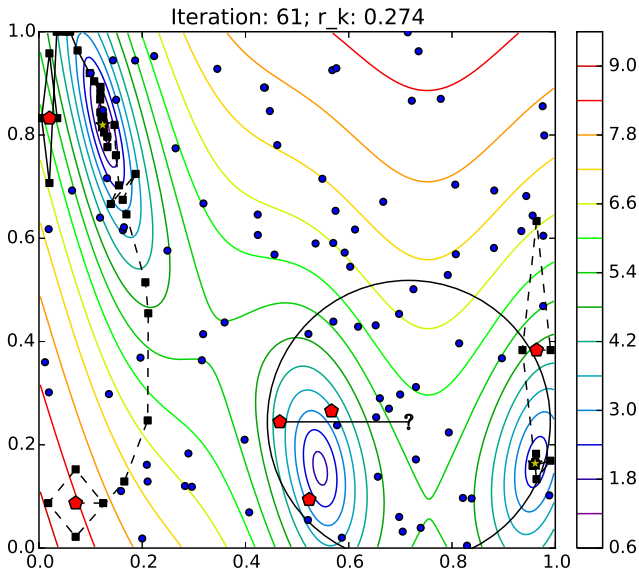
Pausing runs



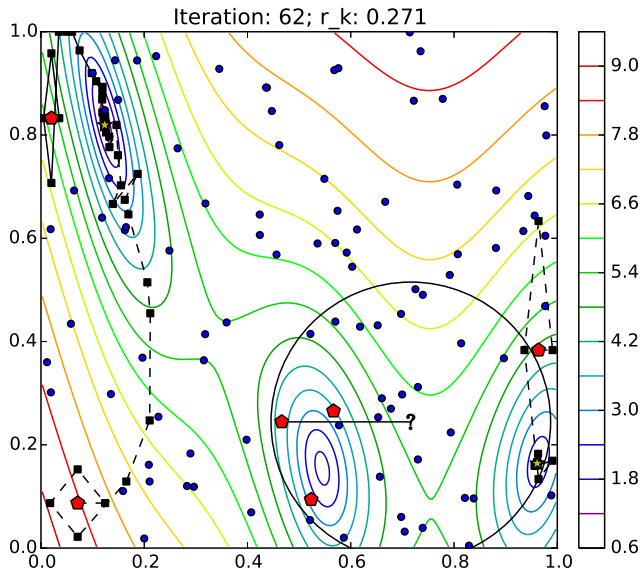
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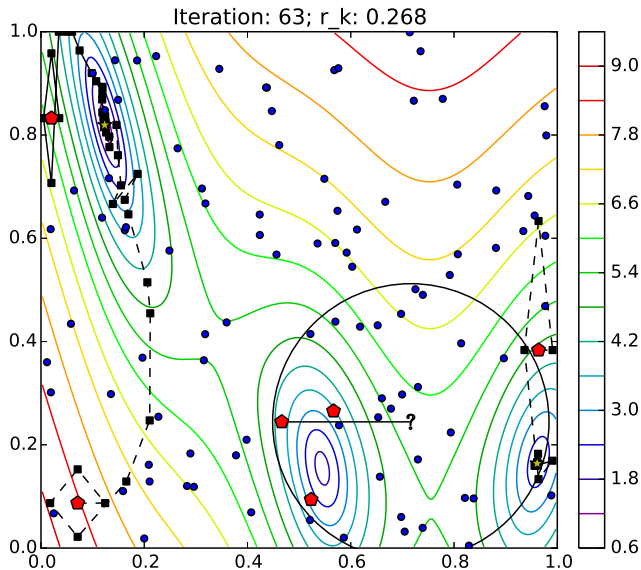
Pausing runs



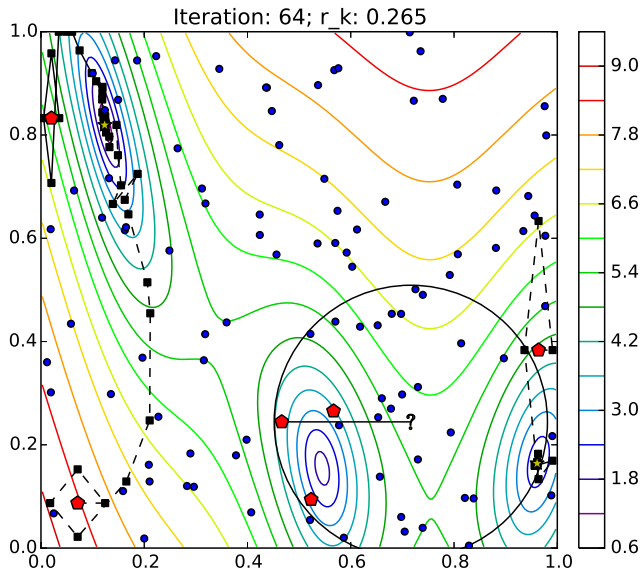
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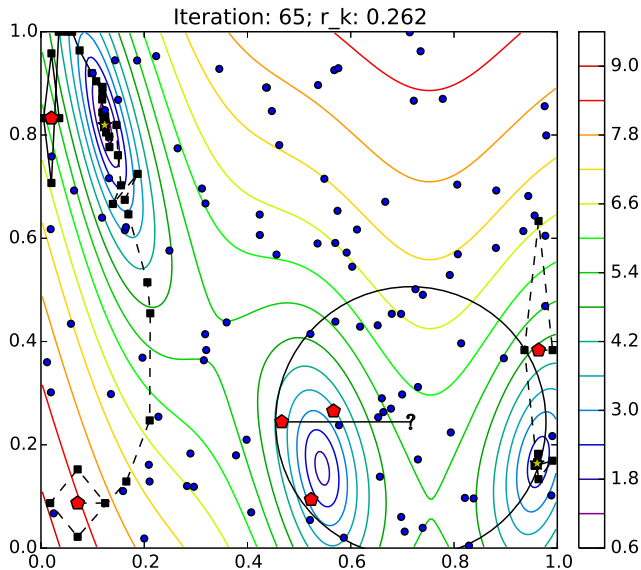
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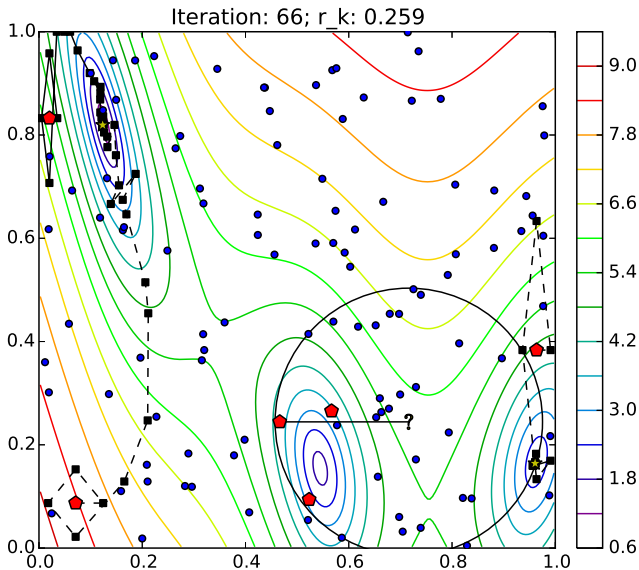
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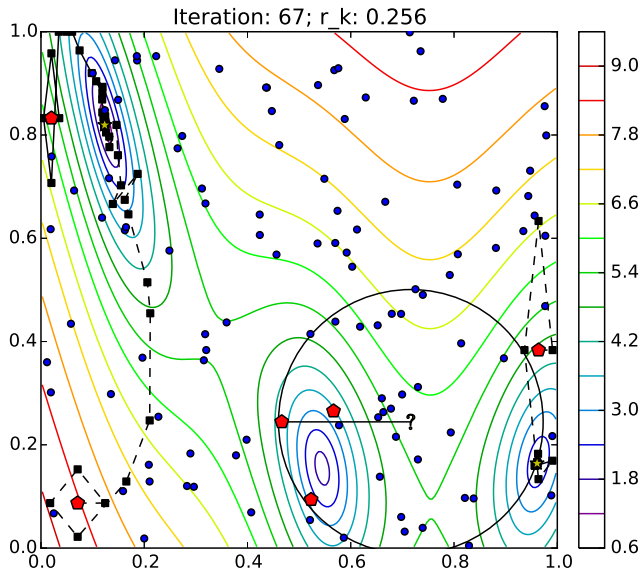
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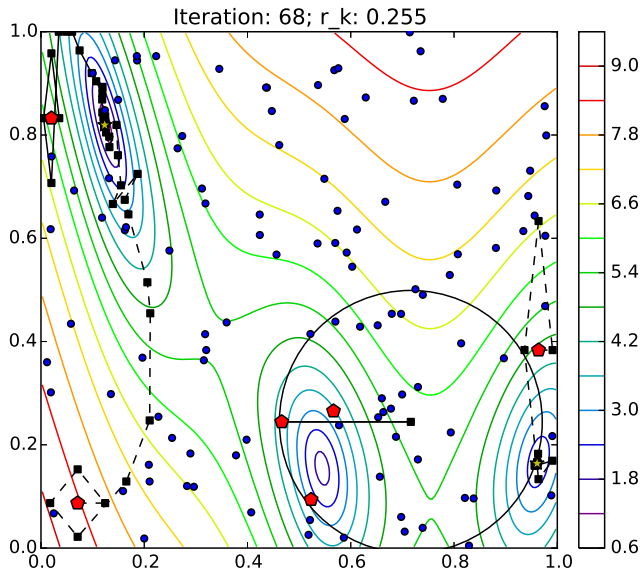
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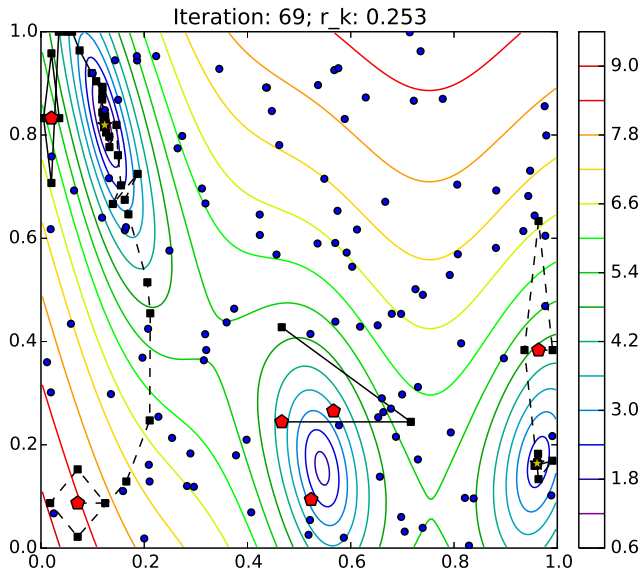
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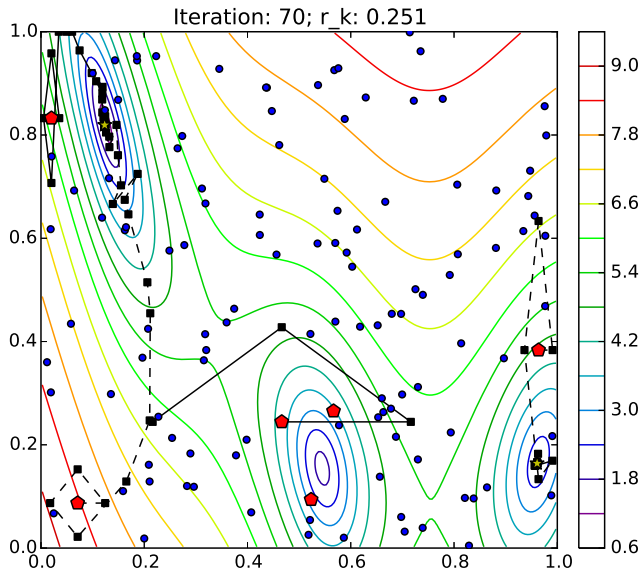
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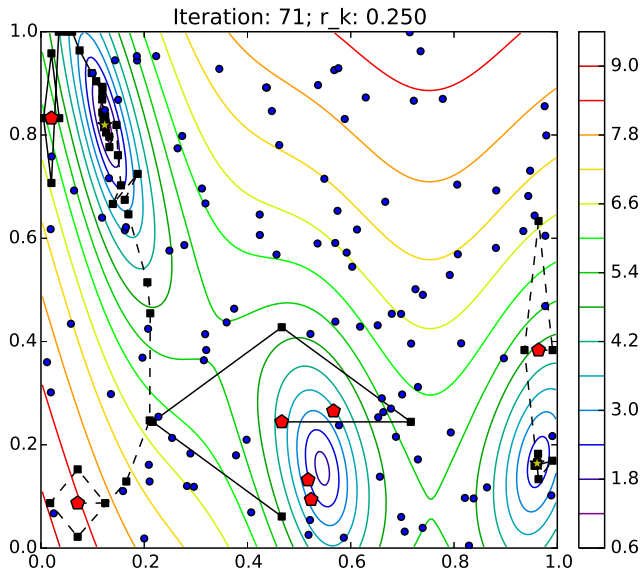
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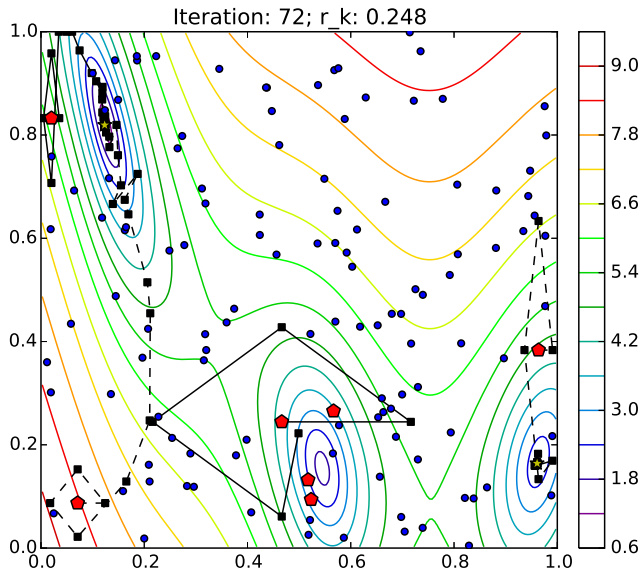
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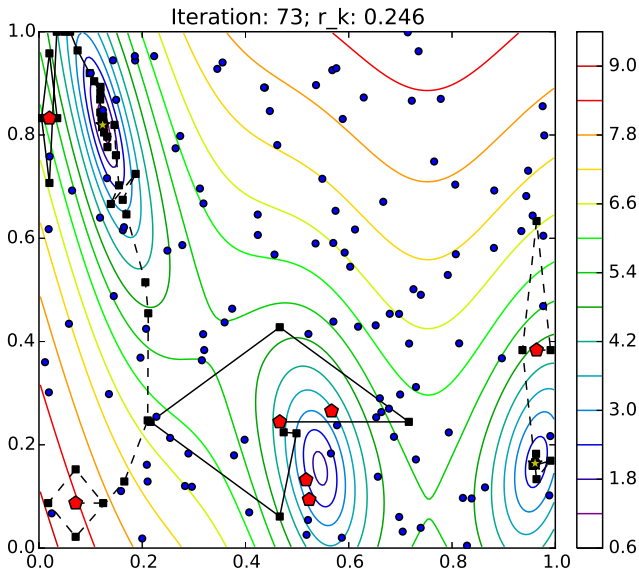
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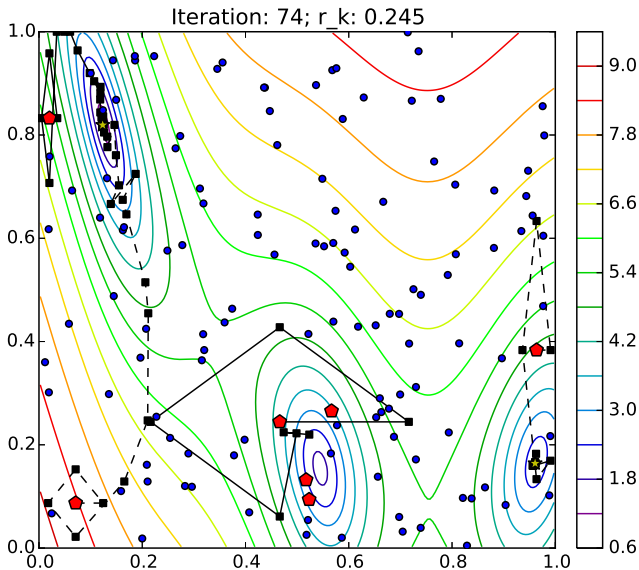
Pausing runs



Pausing runs



Pausing runs



Pausing runs

